

Tensor Decompositions on Emerging Manycore Hardware with Genten and Kokkos

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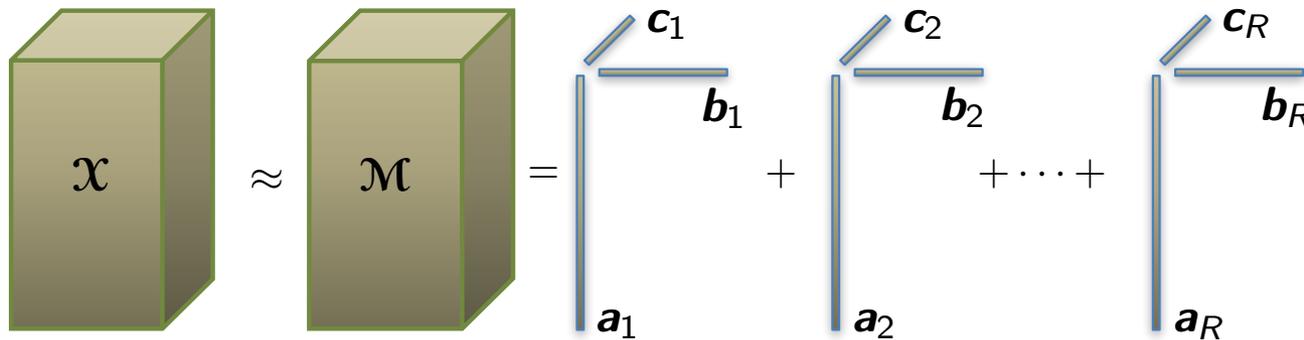
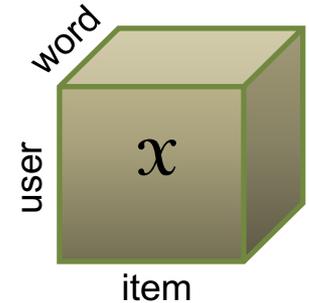
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Genten Project

- Work funded by LDRD program at Sandia National Labs
 - SNL: T. Kolda (PI), K. Devine, E. Phipps, D. Dunlavy, J. Duersch, C. Anderson-Bergman
 - Georgia Tech: Rich Vuduc, Jeff Young, Srinivas Eswar, Jiajia Li
 - Wake Forest: G. Ballard
- Goals of the project:
 - Develop sparse tensor decomposition methods for more general tensor data types (Tammy's talk earlier)
 - Develop distributed memory parallel implementations (Karen's talk)
 - Develop shared memory parallel implementations for next-generation computing architectures (this talk)

Tensors

- N-way array used to represent multi-relationship data
 - E.g., word frequencies in Amazon product rankings:
- Canonical Polyadic (CP) tensor decomposition
 - Approximate tensor as a sum of rank-1 tensors
 - Discovers dominant relationships in data



$$\mathcal{X} \approx \mathcal{M} = \mathbf{a}_1 \circ \mathbf{b}_1 \circ \mathbf{c}_1 + \mathbf{a}_2 \circ \mathbf{b}_2 \circ \mathbf{c}_2 \cdots + \mathbf{a}_R \circ \mathbf{b}_R \circ \mathbf{c}_R$$

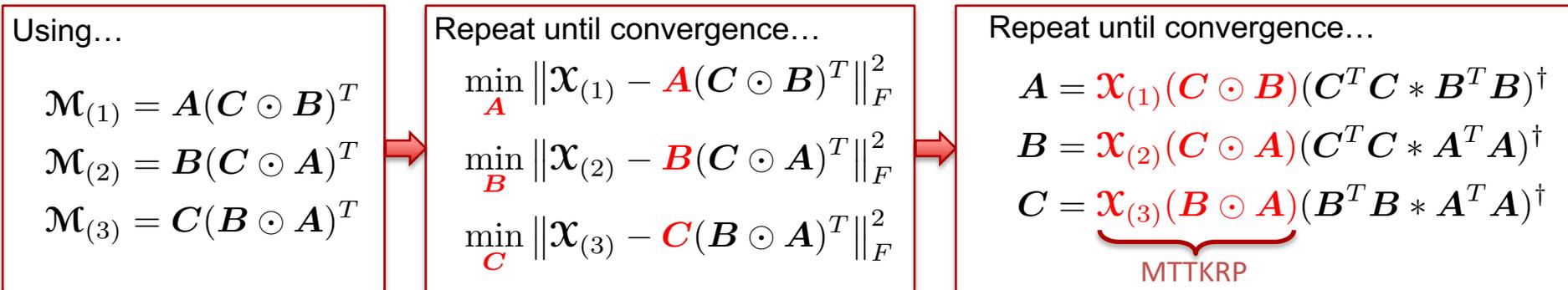
$$x(i, j, k) \approx m(i, j, k) = \sum_{l=1}^R a(i, l)b(j, l)c(k, l)$$

CP via Alternating Least-Squares (ALS)

- CP minimization problem
 - This assumes *Gaussian model* for tensor, resulting in least-squares

$$\min_{\mathcal{M}} \|\mathcal{X} - \mathcal{M}\|_F^2 \quad \text{s.t.} \quad \mathcal{M} = \mathbf{a}_1 \circ \mathbf{b}_1 \circ \mathbf{c}_1 + \cdots + \mathbf{a}_R \circ \mathbf{b}_R \circ \mathbf{c}_R$$

- Solve via alternating linear least squares
 - Fix all but one term, solve linear least squares problem, iterate

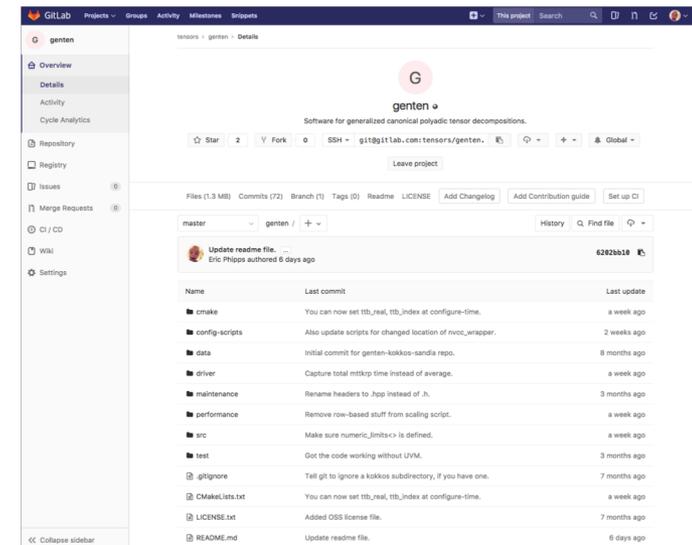


- Matricized Tensor Times Khatri-Rao Product (MTTKRP):
 - Most expensive part of CP-ALS

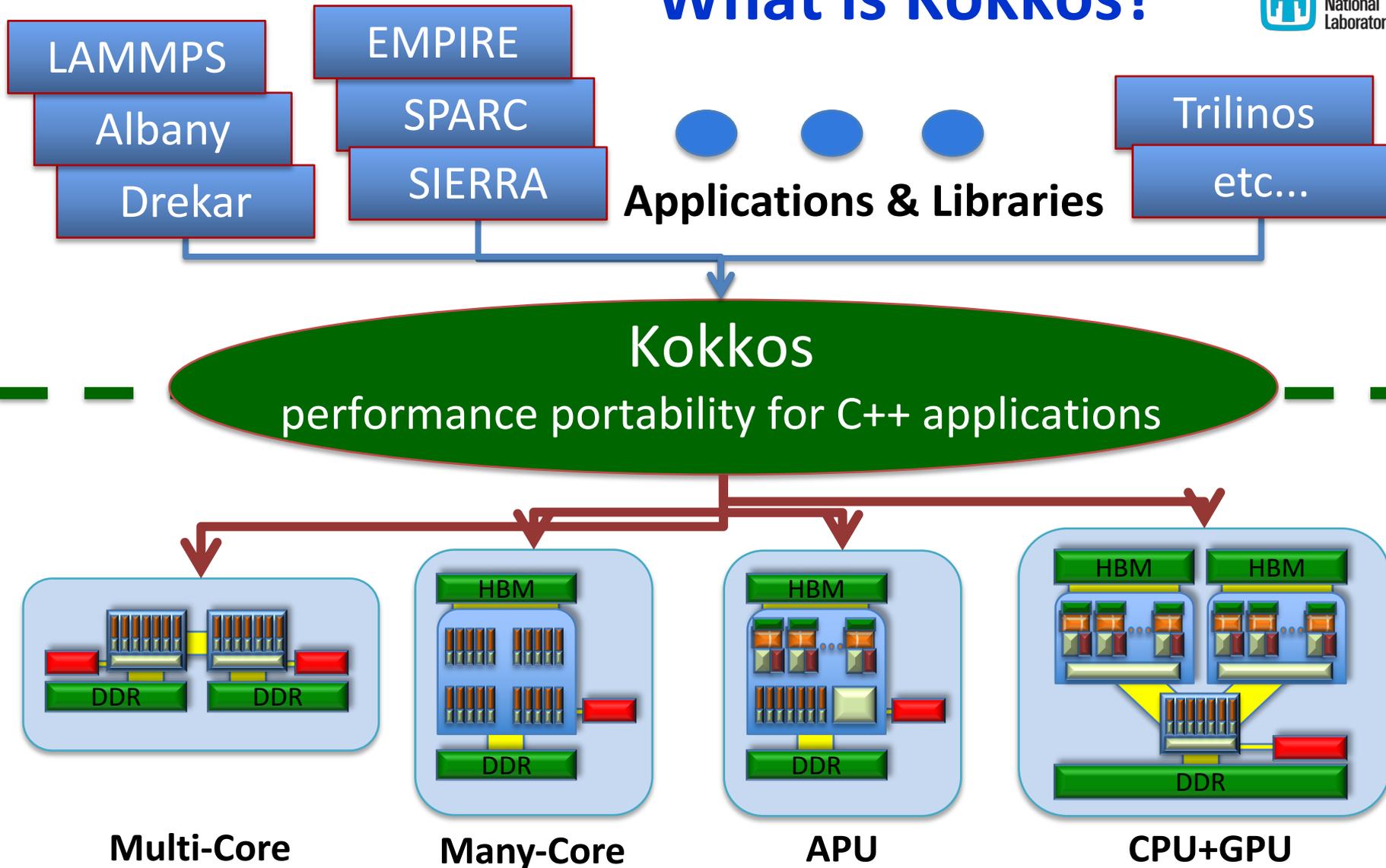
$$\mathbf{V} = \mathcal{X}_{(3)}(\mathbf{A} \odot \mathbf{B}) \implies v(i, l) = \sum_{(j, k) \in \mathcal{N}(\mathcal{X})} x(i, j, k) a(j, l) b(k, l), \quad l = 1, \dots, R$$

Genten : Software for Generalized Canonical Polyadic Tensor Decompositions

- New software package Genten developed at SNL
 - E. Phipps, T. Kolda, D. Dunlavy, G. Ballard, T. Plantenga
 - Based on C++ port of Matlab Tensor Toolbox
 - Publicly available at <https://gitlab.com/tensors/genten>
 - Implements full CP-ALS algorithm for sparse tensors
- Incorporates shared memory parallelism for emerging manycore hardware using Kokkos
 - Multicore CPUs via OpenMP, pThreads
 - GPUs via Nvidia Cuda and AMD ROCm
 - Intel Xeon Phi (a.k.a. KNC/KNL) via OpenMP
- Implements parallelism for all performance-critical operations
 - MTTKRP, tensor inner product, norms, ...
 - Can use optimized third-party libraries (MKL, cuBLAS, ...)
 - Natively handles data transfers between CPU, GPU memory



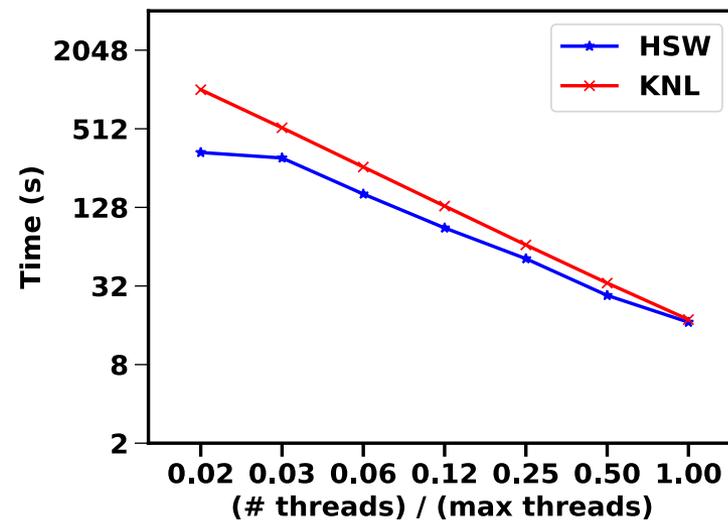
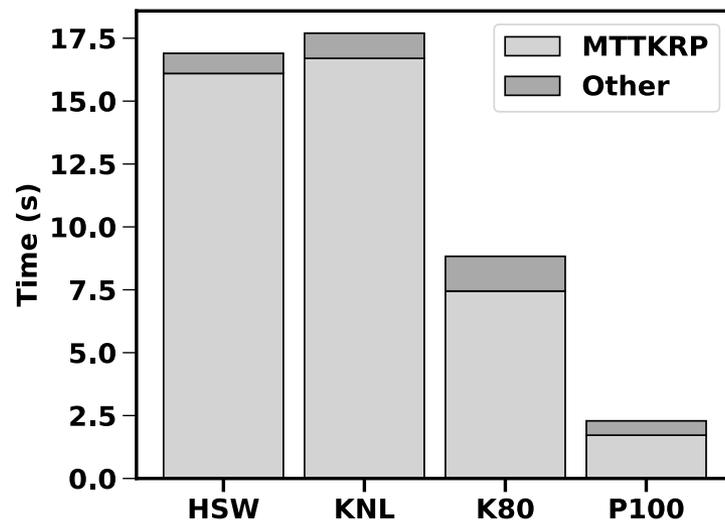
What is Kokkos?



H.C. Edwards, C. Trott, *et al.*, <https://github.com/kokkos/kokkos>

Genten CP-ALS Performance*

- Synthetic data tensor (double precision)
 - 30K x 40K x 50K, 10M non-zeros, R = 128 (perf_CpAlsRandomKtensor performance test in Genten)
 - 10 CP-ALS iterations
- Architectures
 - HSW (OpenMP): Intel Xeon E5-2698v3 CPU, 2.3 GHz, 2 sockets, 16 cores/socket, 2 threads/core
 - KNL (OpenMP): Intel Xeon Phi 7250, 68 cores (using up to 64), 4 threads/core, 16 GB HBM in cache mode
 - K80 (Cuda): Nvidia Kepler K80 GPU, 12 GB GDDR5
 - P100 (Cuda): Nvidia Pascal P100 GPU, 16 GB HBM



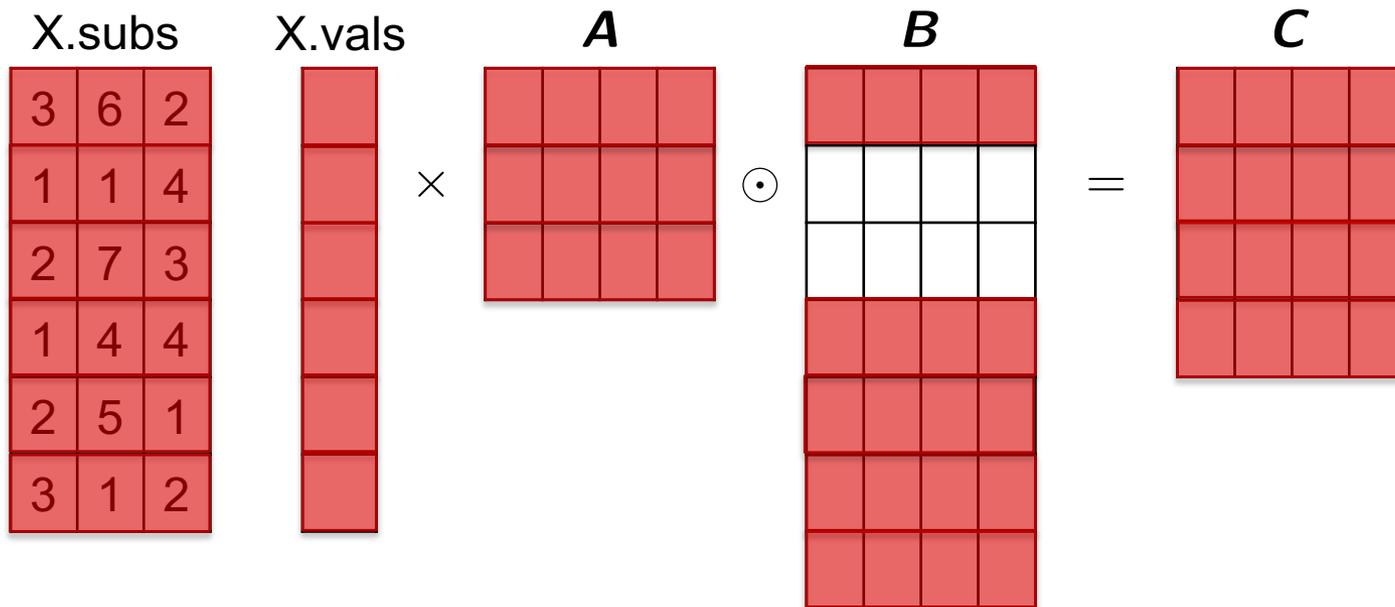
*E. Phipps and T. Kolda, *Software for Sparse Tensor Decomposition on Emerging Computing Architectures*, submitted to SIAM SISC, [arXiv.org:1809.09175](https://arxiv.org/abs/1809.09175).

MTTKRP

- Sparse tensors stored in coordinate (COO) format
 - List of nonzero values and coordinates (subscripts)

$$\mathcal{X} \in \mathbb{R}^{3 \times 7 \times 4}, \quad R = 4$$

$$C = \mathcal{X}_{(3)}(A \odot B) \implies c(k, l) = \sum_{(i, j, k) \in \mathcal{N}(\mathcal{X})} x(i, j, k) a(i, l) b(j, l), \quad l = 1, \dots, R$$



Serial MTTKRP Algorithm

```
Sptensor X;  
Ktensor u;  
FacMatrix v;  
  
const int nnz = X.nnz();  
for (int i=0; i<nnz; ++i)  
{  
    double *tmp = new double[R];  
  
    const size_t k = X.subscript(i,n);  
    const double x_val = X.value(i);  
    for (int j = 0; j < R; ++j)  
        tmp[j] = x_val;  
  
    for (int m = 0; m < N; m++) {  
        if (m != n) {  
            const double *row = u[m].rowptr(X.subscript(i,m));  
            for (int j = 0; j < R; ++j)  
                tmp[j] *= row[j];  
        }  
    }  
  
    for (int j = 0; j < R; ++j)  
        v.entry(k,j) += tmp[j];  
  
    delete [] tmp;  
}
```

Matricized Tensor Times Khatri-Rao Product (MTTKRP)

- Factor matrices stored row-wise (Kokkos::LayoutRight)

$$\mathbf{V} = \mathbf{x}_{(n)} (\mathbf{U}^N \odot \dots \odot \mathbf{U}^{n+1} \odot \mathbf{U}^{n-1} \odot \dots \odot \mathbf{U}^1)$$
$$v(i_n, :) = \sum_{i \in \mathcal{N}(\mathbf{X})} x_i u^1(i_1, :) \dots u^{n-1}(i_{n-1}, :) u^{n+1}(i_{n+1}, :) \dots u^N(i_N, :)$$

Flat Parallelism Over Tensor Non-zeros

```
typedef Kokkos::RangePolicy<ExecSpace> Policy;
const int nnz = X.nnz();
Policy policy(0,nnz);
Kokkos::parallel_for( policy,
                    KOKKOS_LAMBDA(const int i)
{
    double *tmp = new double[R];

    const size_t k = X.subscript(i,n);
    const double x_val = X.value(i);

    for (int j = 0; j < R; ++j)
        tmp[j] = x_val;

    for (int m = 0; m < N; m++) {
        if (m != n) {
            const double *row = u[m].rowptr(X.subscript(i,m));
            for (int j = 0; j < R; ++j)
                tmp[j] *= row[j];
        }
    }

    for (int j = 0; j < R; ++j)
        Kokkos::atomic_add(&v.entry(k,j), tmp[j]);

    delete [] tmp;
});
```



Parallelize MTTKRP over tensor non-zeros



Multiple threads writing to same row of V requires atomic update (not the only way to do it)

Flat Parallelism Over Tensor Non-zeros

```
typedef Kokkos::RangePolicy<ExecSpace> Policy;
const int nnz = X.nnz();
Policy policy(0,nnz);
Kokkos::parallel_for( policy,
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{
    double *tmp = new double[R];

    const size_t k = X.subscript(i,n);
    const double x_val = X.value(i);

    for (int j = 0; j < R; ++j)
        tmp[j] = x_val;

    for (int m = 0; m < N; m++) {
        if (m != n) {
            const double *row = u[m].rowptr(X.subscript(i,m));
            for (int j = 0; j < R; ++j)
                tmp[j] *= row[j];
        }
    }

    for (int j = 0; j < R; ++j)
        Kokkos::atomic_add(&v.entry(k,j), tmp[j]);

    delete [] tmp;
});
```

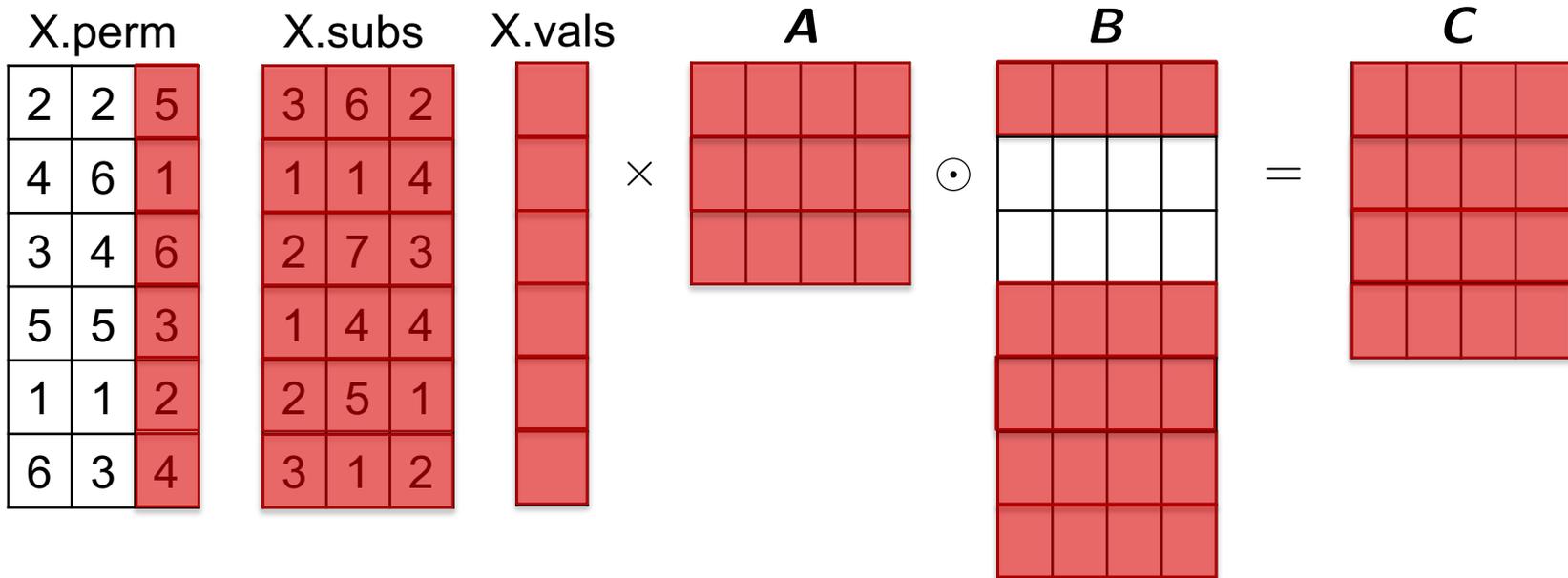
Improvements in MTTKRP implementation:

- Thread teams to process in parallel:
 - Tensor non-zeros
 - Factor matrix columns
- Thread-scalable mechanism for tmp buffer
 - Kokkos scratch-pad memory (MTTKRP-A)
 - Thread-local, compile-time polymorphic arrays (MTTKRP-B)
- Blocking on R for better cache performance, occupancy
- Experimenting with alternatives to atomics
 - Thread privatization buffers

Permuted MTTKRP Inspired by COO

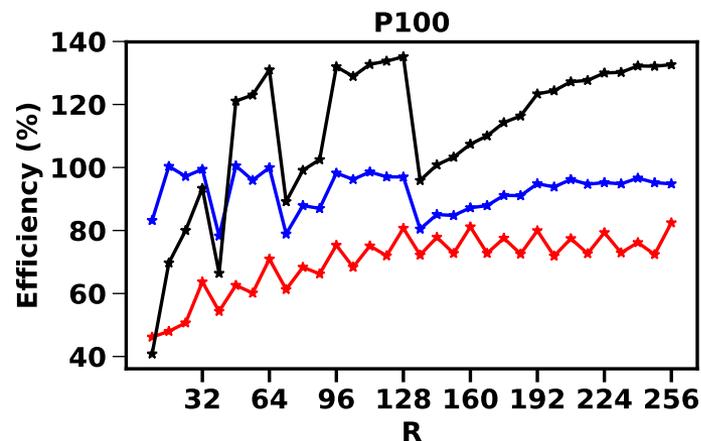
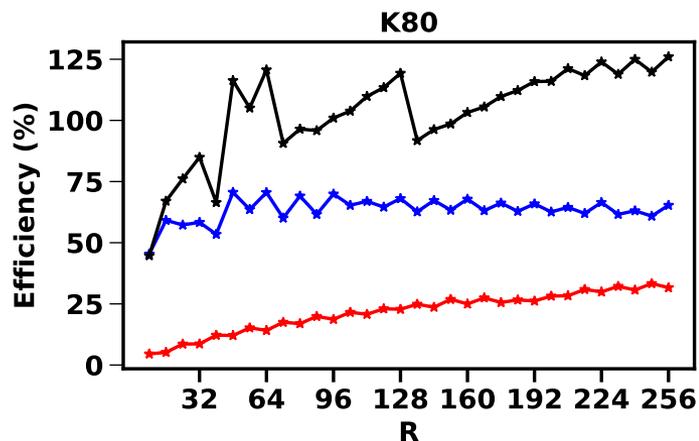
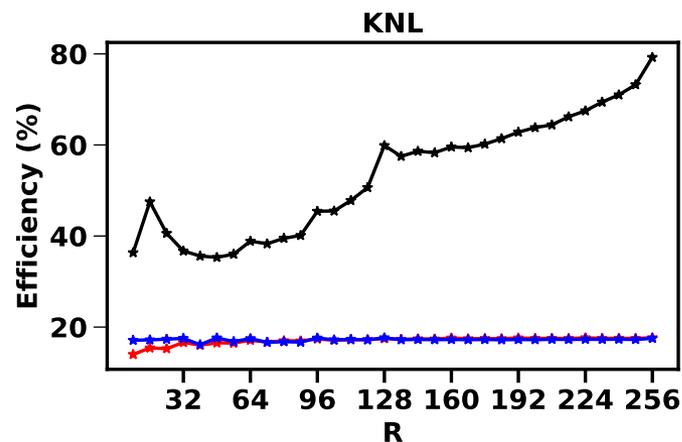
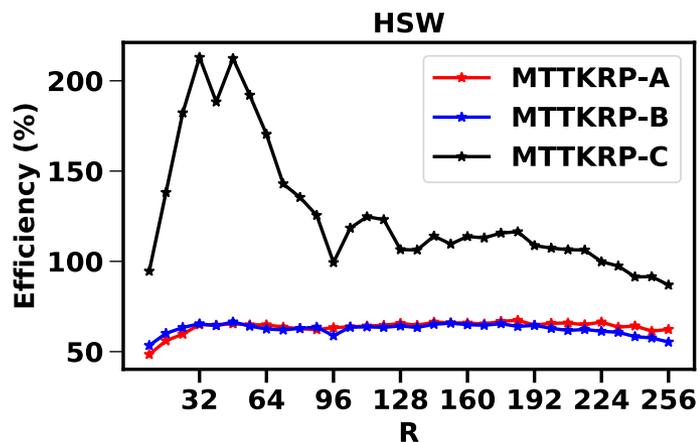
Mat-Vec

- Poorly structured tensors result in high atomic contention
 - COO sparse matrix-vector product leverages sorting by increasing row-index
- Apply this to coordinate-based sparse tensor format
 - Don't want multiple copies of tensor with different sortings
 - Instead, compute permutation array p so that $X.\text{subscript}(p(:,n),n)$ is increasing
 - Each thread only writes when row index $X.\text{subscript}(p(:,n),n)$ changes
 - Substantially reduces number of atomic updates and improves their locality



MTTKRP Percentage of Peak Bandwidth

$$\text{GB/s} = \left((NR + 3) \times \text{sizeof}(\text{Real}) + N \times \text{sizeof}(\text{Ordinal}) \right) \times \frac{NNZ}{t}$$



- MTTKRP-A: Kokkos scratch-pad arrays
- MTTKRP-B: Compile-time polymorphic arrays

- MTTKRP-C: Permutation-based MTTKRP

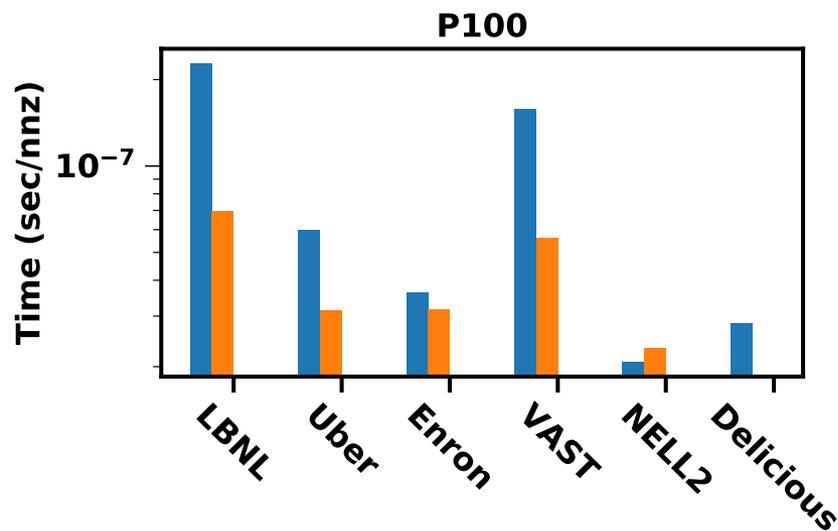
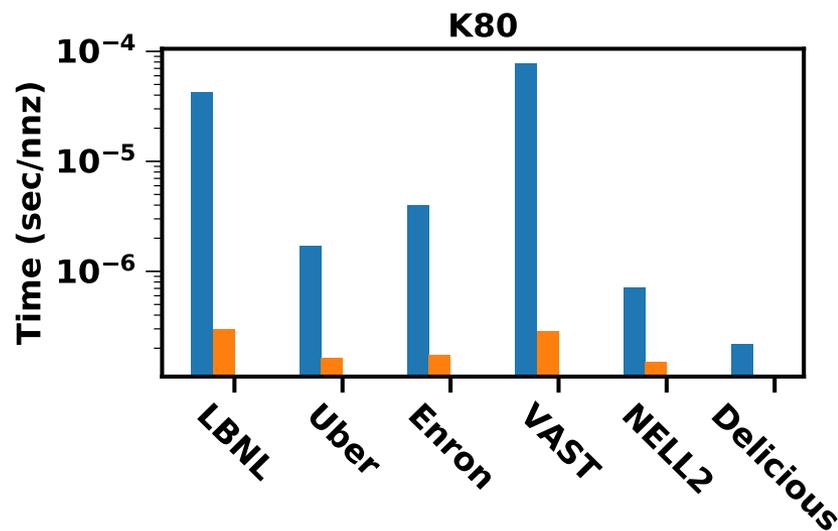
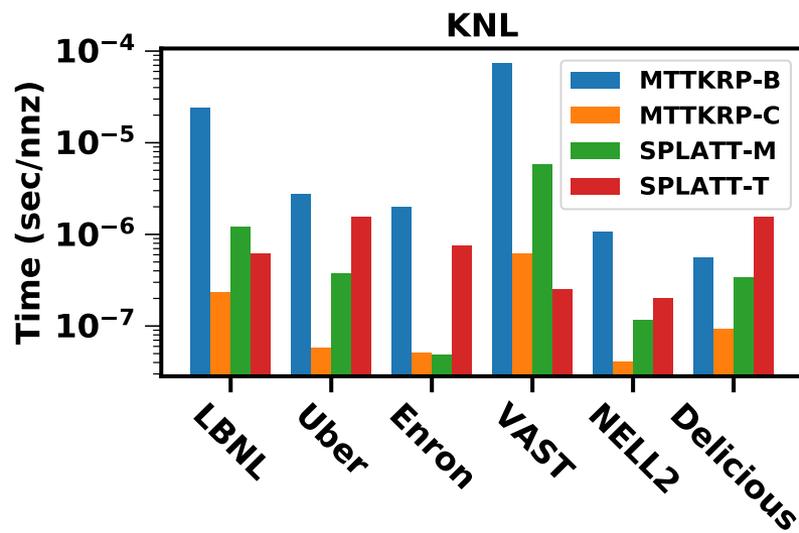
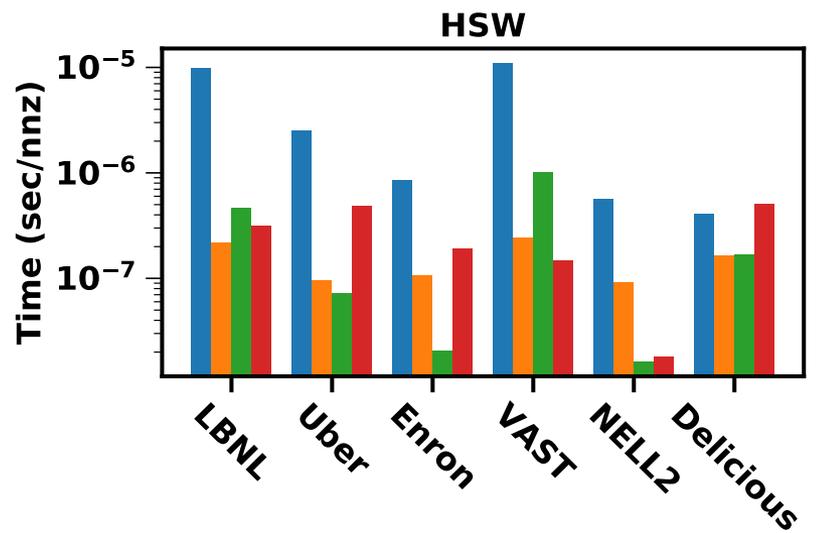
Performance on Real Tensors

- Look at Genten MTTKRP performance on several tensors from the FROSTT collection
 - Shaden Smith et al, <https://frost.io>
 - Fixed 10 iterations of CP-ALS
 - $R = 16$ for all tests

Name	Order	Dimensions	Nonzeros
LBNL	5	$1.6\text{K} \times 4.2\text{K} \times 1.6\text{K} \times 4.2\text{K} \times 868\text{K}$	1.7M
Uber	4	$183 \times 24 \times 1.1\text{K} \times 1.7\text{K}$	13M
Enron	4	$6.0\text{K} \times 5.7\text{K} \times 244\text{K} \times 1.2\text{K}$	54M
VAST	5	$165\text{K} \times 11\text{K} \times 2 \times 100 \times 89$	26M
NELL2	3	$12\text{K} \times 9.1\text{K} \times 29\text{K}$	77M
Delicious	4	$532\text{K} \times 17\text{M} \times 2.5\text{M} \times 1.4\text{K}$	140M

- Compare MTTKRP-B/C to SPLATT on OpenMP architectures
 - Shaden Smith et al, <http://glaros.dtc.umn.edu/gkhome/splatt/overview>
 - Compressed Sparse Fiber (CSF) storage format (Smith and Karypis, 2015)
 - Use default 2 CSF modes, with and without tiling

Total MTTKRP Time per Nonzero



Summary & Conclusions

- New software package Genten
 - Publicly available at <https://gitlab.com/tensors/genten>
 - POC: Eric Phipps, etphipp@sandia.gov
- Built on Kokkos for shared-memory parallelism and high performance, portable to emerging parallel architectures
- Simple modification of standard MTTKRP algorithm to store and iterate over permutation arrays improves performance
- Future Plans
 - Continue to improve performance (MTTKRP, sorting, ...)
 - Kokkos implementation of state-of-the-art formats such as CSF and HiCOO
 - Incorporate generalized tensor decomposition R&D by Kolda et al for general tensor data types
 - Incorporate distributed memory parallelism (K. Devine)

Acknowledgments

- Thanks to T. Kolda and others for many helpful comments!
- This work was supported by the Laboratory Directed Research and Development program at Sandia National Laboratories, a multimission laboratory managed and operated by National Technology and Engineering Solutions of Sandia, LLC., a wholly owned subsidiary of Honeywell International, Inc., for the U.S. Department of Energy's National Nuclear Security Administration under contract DE-NA-0003525.

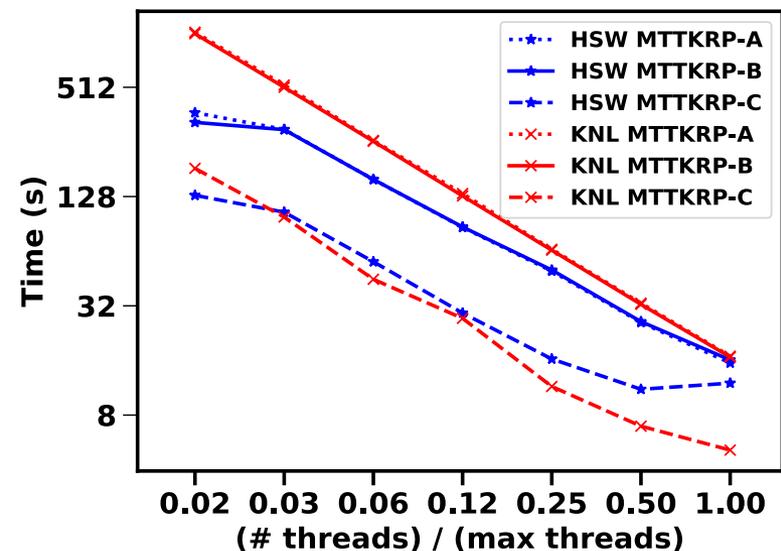
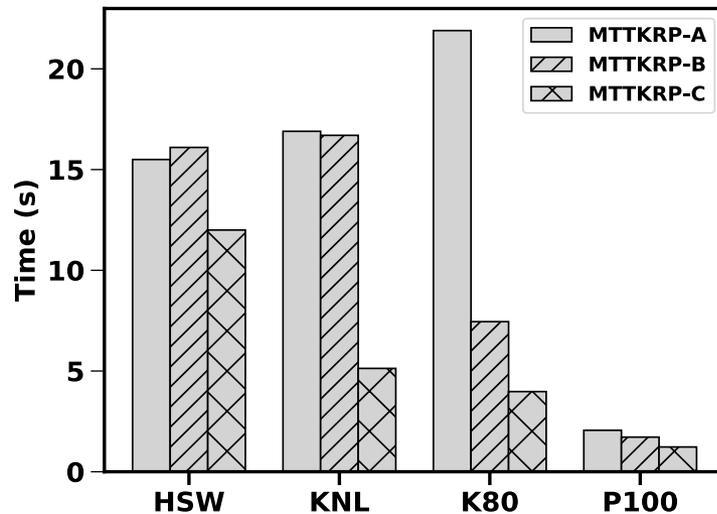
Backup Slides

Improving MTTKRP Performance

- Problems with current algorithm
 - High atomic contention if tensor is poorly structured
 - Poor-locality in writes to rows of result factor matrix
- Take inspiration from coordinate-based (COO) sparse mat-vec
 - Nonzero indices are sorted by increasing row index
 - Allows marching down the rows, improving locality
 - Reduce partial sums across non-zeros with the same row index to reduce atomic writes
- Apply this to coordinate-based sparse tensor format
 - For each mode n , compute permutation array p so that $X.\text{subscript}(p(:,n),n)$ is increasing
 - Requires N sorting operations at setup (cost of a few MTTKRP)
 - Store 2-D permutation array same size as coordinates ($\sim 2x$ increased memory usage)
 - MTTKRP algorithm:
 - Each thread team loops over blocks of non-zeros in $p(\cdot)$ order
 - Accumulate result until change in row index
 - Atomic write if first or last row in block, otherwise normal write
- *Substantially reduces number of atomic updates and improves their locality*

Permuted MTTKRP Performance

- MTTKRP-A: Team-based MTTKRP algorithm w/Kokkos scratch-pad arrays
 - Store tmp buffer in shared memory on GPU
- MTTKRP-B: Team-based MTTKRP algorithm w/compile-time polymorphic arrays
 - Store tmp buffer in registers on GPU
- MTTKRP-C: Team-based MTTKRP w/permutation arrays and compile-time polymorphic arrays for tmp buffers



Total MTTKRP Time

Name	Order	Dimensions	Nonzeros
LBNL	5	1.6K × 4.2K × 1.6K × 4.2K × 868K	1.7M
Uber	4	183 × 24 × 1.1K × 1.7K	13M
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VAST	5	165K × 11K × 2 × 100 × 89	26M
NELL2	3	12K × 9.1K × 29K	77M
Delicious	4	532K × 17M × 2.5M × 1.4K	140M

- MTTKRP-B: MTTKRP w/compile-time polymorphic arrays
- MTTKRP-C: Permutation-based MTTKRP

Arch.	Method	LBNL	Uber	Enron	VAST	NELL2	Delicious
HSW	MTTKRP-B	16.6	8.3	45.7	284.0	42.8	57.4
	MTTKRP-C	0.4	0.3	5.7	6.3	7.0	22.7
	SPLATT-M	0.8	0.2	1.1	26.1	1.3	23.4
	SPLATT-T	0.5	1.6	10.4	3.8	1.4	71.2
KNL	MTTKRP-B	41.2	9.2	108.0	1910.0	82.4	79.1
	MTTKRP-C	0.4	0.2	2.8	16.1	3.2	12.9
	SPLATT-M	2.0	1.2	2.6	150.5	9.0	47.7
	SPLATT-T	1.0	5.1	41.3	6.5	15.6	219.1
K80	MTTKRP-B	72.4	5.6	218.0	2010.0	55.2	30.3
	MTTKRP-C	0.5	0.5	9.4	7.4	11.6	–
P100	MTTKRP-B	0.4	0.2	2.0	4.1	1.6	4.0
	MTTKRP-C	0.1	0.1	1.7	1.5	1.8	–

Tensor Sorting Cost

- Permutation method requires N sorting operations
 - Sort tensor coordinates for each mode
 - Small, but not inconsequential cost
- Sorting methods used
 - OpenMP: Parallel stable sort from Intel
 - Cuda: Thrust stable sort

Sorting time divided by avg. (permuted) CP-ALS iteration time

Architecture	LBNL	Uber	Enron	VAST	NELL2	Delicious
HSW	0.6	4.3	9.2	5.3	6.2	4.0
KNL	1.0	5.5	8.6	1.3	7.6	4.9
K80	1.1	3.0	3.3	3.3	2.6	–
P100	0.5	1.3	4.2	3.6	4.6	–

CP via Alternating Least-Squares (ALS)

*For least-squares error model (Gaussian approximation),
solve minimization by fixing all but one term, then iterate*

Input: \mathbf{X} , R , $MaxIts$, tol

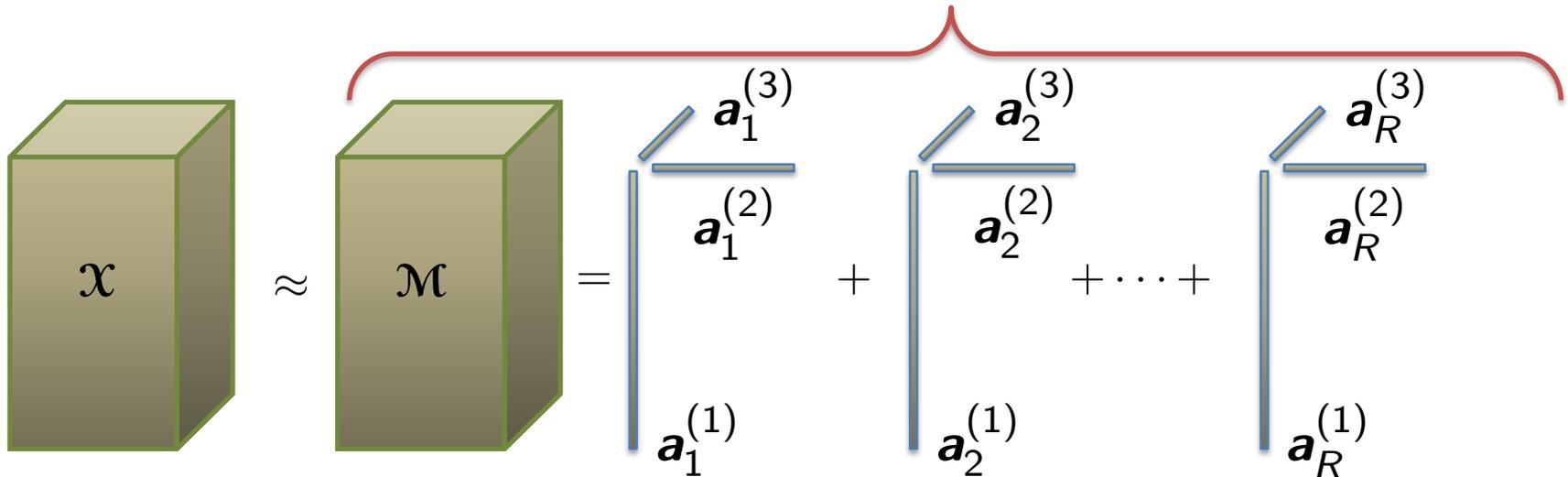
- 1: Initialize $\mathbf{A}^{(1)}, \dots, \mathbf{A}^{(N)}$
 - 2: $i \leftarrow 0$
 - 3: **repeat**
 - 4: **for** $n = 1, \dots, N$ **do**
 - 5: $\mathbf{V}^{(n)} \leftarrow (\mathbf{A}^{(1)})^T \mathbf{A}^{(1)} * \dots * (\mathbf{A}^{(n-1)})^T \mathbf{A}^{(n-1)} * (\mathbf{A}^{(n+1)})^T \mathbf{A}^{(n+1)} * \dots * (\mathbf{A}^{(N)})^T \mathbf{A}^{(N)}$ // Gramian
 - 6: $\mathbf{B}^{(n)} \leftarrow \mathbf{X}_{(n)} (\mathbf{A}^{(N)} \odot \dots \odot \mathbf{A}^{(n+1)} \odot \mathbf{A}^{(n-1)} \odot \dots \odot \mathbf{A}^{(1)})$ // MTTKRP
 - 7: $\mathbf{A}^{(n)} \leftarrow \mathbf{B}^{(n)} (\mathbf{V}^{(n)})^\dagger$ // Solve
 - 8: $\lambda_k \leftarrow \|\mathbf{a}_k^{(n)}\|$, $k = 1, \dots, R$ // ColumnNorm
 - 9: $\mathbf{A}^{(n)} \leftarrow \mathbf{A}^{(n)} \mathbf{\Lambda}^{-1}$ // ColumnScale
 - 10: **end for**
 - 11: $i \leftarrow i + 1$
 - 12: **until** $i \geq MaxIts$ **or** change in $\|\mathbf{X} - [\mathbf{\lambda}; \mathbf{A}^{(1)}, \dots, \mathbf{A}^{(N)}]\| < tol$
- Output:** $\mathbf{\lambda}$, $\mathbf{A}^{(1)}, \dots, \mathbf{A}^{(N)}$

Canonical Polyadic (CP) Decomposition

Approximate tensor as a sum of rank-1 tensors to find dominant relationships

Data

CP Model



$$\mathbf{x}_{i_1, i_2, \dots, i_N} \approx \mathcal{M}_{i_1, i_2, \dots, i_N} = \sum_{j=1}^R \lambda_j \mathbf{a}_{i_1, j}^{(1)} \mathbf{a}_{i_2, j}^{(2)} \cdots \mathbf{a}_{i_N, j}^{(N)}$$

$$\begin{aligned} \min_{\mathcal{M}} \quad & \|\mathcal{X} - \mathcal{M}\| \\ \text{s.t.} \quad & \mathcal{M} = \llbracket \boldsymbol{\lambda}; \mathbf{A}^{(1)}, \mathbf{A}^{(2)}, \dots, \mathbf{A}^{(N)} \rrbracket, \quad \mathbf{A}^{(n)} = [\mathbf{a}_1^{(n)}, \dots, \mathbf{a}_R^{(n)}] \end{aligned}$$

What is Kokkos?

- **ΚÓΚΚΟΣ** (Greek, not an acronym)
 - Translation: “granule” or “grain” ; *like grains of sand on a beach*
 - H.C. Edwards, C. Trott, *et al*
- Performance Portable Thread-Parallel Programming Model
 - E.g., “X” in “MPI+X” ; not a distributed-memory programming model
 - Application identifies its parallelizable grains of computations *and* data
 - Kokkos maps those computations onto cores *and* that data onto memory
- Fully Performance Portable C++11 Library Implementation
 - *Not* a language extension (e.g., OpenMP, OpenACC, OpenCL, ...)
 - Production – open source at <https://github.com/kokkos/kokkos>
 - ✓ Compilers: GNU, LLVM, Intel, NVIDIA, IBM XL, Cray
 - ✓ Multicore CPU - including NUMA architectural concerns
 - ✓ Intel Xeon Phi (KNC/KNL)
 - ✓ NVIDIA GPU (Kepler, Pascal, Volta)
 - ✓ IBM Power 8/9
 - AMD Fusion

Advantages of Kokkos (my view)

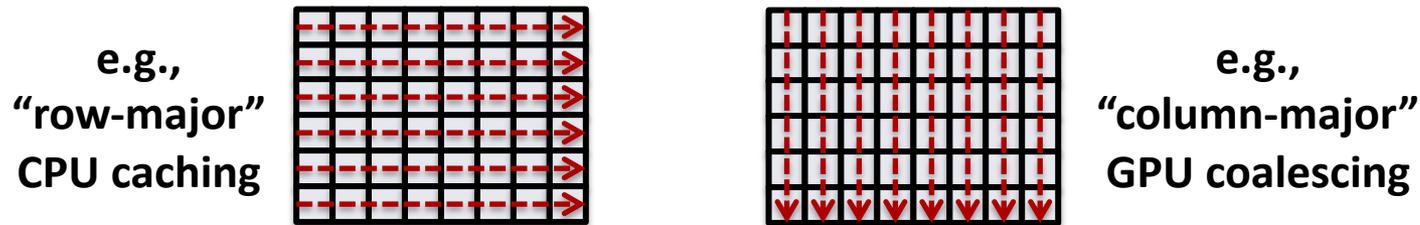
- Tools to build portable, thread-parallel algorithms with good performance
 - Tools to manage memory allocations and multiple memory spaces
 - Ability to introduce data layout polymorphism
 - Simple, portable interfaces to advanced architecture-specific capabilities
 - Light-weight abstractions that don't impede performance
- Future-proofing
 - Substantial buy-in from DOE and vendors
 - Expected to be part of C++-2x standard
- Flexibility to write directly to the underlying programming model
 - Can always write specializations specific to an architecture if necessary/desired

Kokkos Abstractions

- Parallel Pattern of user's computations
 - parallel_for, parallel_reduce, parallel_scan, task-graph, ... (*extensible*)
- Execution Policy tells *how* user computation will be executed
 - Static scheduling, dynamic scheduling, thread-teams, ... (*extensible*)
- Execution Space tells *where* user computations will execute
 - Which cores, numa region, GPU, ... (*extensible*)
- Memory Space tells *where* user data resides
 - Host memory, GPU memory, high bandwidth memory, ... (*extensible*)
- Layout (policy) tells *how* user data is laid out in memory
 - Row-major, column-major, array-of-struct, struct-of-array ... (*extensible*)
- Differentiating: Layout and Memory Space
 - Versus other programming models (OpenMP, OpenACC, ...)
 - Critical for performance portability ...

Layout Abstraction: Multidimensional Array

- **Classical (50 years!) data pattern for science & engineering codes**
 - Computer languages hard-wire multidimensional array layout mapping
 - Problem: different architectures *require* different layouts for performance
 - **Leads to architecture-specific versions of code to obtain performance**
 - E.g., “Array of Structure” ↔ “Structure of Array” redesigns



■ Kokkos *separates* layout from user’s computational code

- Choose layout for architecture-specific memory access pattern
 - **Without modifying user’s computational code**
- **Polymorphic layout** via C++ template meta-programming (*extensible*)
 - e.g., Hierarchical Tiling layout (array of structure of array)
- **Bonus: easy/transparent use of special data access hardware**
 - Atomic operations, GPU texture cache, ... (*extensible*)

Integrated mapping of users' parallel computations *and* data through abstractions of patterns, policies, spaces, *and* layout.

- **Versus other thread parallel programming models (mechanisms)**
 - OpenMP, OpenACC, OpenCL, ... have parallel execution
 - OpenMP 4 finally has execution spaces; when memory spaces ??
 - **All of these neglect data layout mapping**
 - Requiring significant code refactoring to change data access patterns
 - Cannot provide *performance* portability
 - **All require language and compiler changes for extension**
- **Kokkos extensibility “future proofing” for evolving architectures**
 - Library extensions, not compiler extensions
 - E.g., Intel KNL high bandwidth memory ← just another memory space
- **Productivity versus other programming models?**

Flat Parallelism Over Tensor Non-zeros

```
typedef Kokkos::RangePolicy<ExecSpace> Policy;
const int nnz = X.nnz();
Policy policy(0,nnz);
Kokkos::parallel_for( policy,
                    KOKKOS_LAMBDA(const int i)
{
    double *tmp = new double[R]; ← Dynamic memory allocation (very bad on GPU)

    const size_t k = X.subscript(i,n);
    const double x_val = X.value(i);

    for (int j = 0; j < R; ++j)
        tmp[j] = x_val;

    for (int m = 0; m < N; m++) {
        if (m != n) {
            const double *row = u[m].rowptr(X.subscript(i,m));
            for (int j = 0; j < R; ++j)
                tmp[j] *= row[j]; ← Not exploiting parallelism over columns
        }
    }

    for (int j = 0; j < R; ++j)
        Kokkos::atomic_add(&v.entry(k,j), tmp[j]);

    delete [] tmp;
});
```

Kokkos Hierarchical Parallelism

- Kokkos maps a league of teams across a given index range
 - Each team consists of a given number of “threads”
 - Each “thread” can iterate over a “vector” range
- Team:
 - CPU/Phi: One or more cores
 - GPU: Cuda thread block
- Thread:
 - CPU/Phi: Hyperthread in a core
 - GPU: Warp in a thread block
- Vector:
 - CPU/Phi: Vector lane in a SIMD instruction
 - GPU: Thread within a warp

Adding Scratch Space and Vector Parallelism

```
typedef Kokkos::TeamPolicy<ExecSpace> Policy;
const int nnz = X.nnz();
const int VectorSize = 8;
const int TeamSize = 1;
Policy policy(nnz,TeamSize,VectorSize);

typedef Kokkos::View< double*,
  ExecSpace::scratch_memory_space,
  Kokkos::MemoryUnmanaged > ScratchSpace;
const size_t bytes = ScratchSpace::shmem_size(R);

Kokkos::parallel_for(
  policy.set_scratch_size(0,Kokkos::PerTeam(bytes)),
  KOKKOS_LAMBDA(Policy::member_type team)
{
  const size_t i = team.league_rank();

  ScratchSpace tmp(team.team_scratch(0), R);

  const size_t k = X.subscript(i,n);
  const double x_val = X.value(i);
```

```
Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
  [&] (const int r)
{
  tmp[r] = x_val;
});

for (int m = 0; m < N; ++m) {
  if (m != n) {
    const double *row = u[m].rowptr(X.subscript(i,m));
    Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
      [&] (const int r)
      {
        tmp[r] *= row[r];
      });
  }
}

Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
  [&] (const int r)
{
  Kokkos::atomic_add(&v.entry(k,r), tmp[r]);
});
```



Vector parallelism across columns of factor matrices

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const int VectorSize = 8;
const int TeamSize = 1;
Policy policy(nnz,TeamSize,VectorSize);

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    ExecSpace::scratch_memory_space,
    Kokkos::MemoryUnmanaged > ScratchSpace;
const size_t bytes = ScratchSpace::shmem_size(R);

Kokkos::parallel_for(
    policy.set_scratch_size(0,Kokkos::PerTeam(bytes)),
    KOKKOS_LAMBDA(Policy::member_type team)
{
    const size_t i = team.league_rank();

    ScratchSpace tmp(team.team_scratch(0), R);

    const size_t k = X.subscript(i,n);
    const double x_val = X.value(i);
```



Allocation from scratch memory space
(shared memory on GPU)

```
Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
    [&] (const int r)
{
    tmp[r] = x_val;
});

for (int m = 0; m < N; ++m) {
    if (m != n) {
        const double *row = u[m].rowptr(X.subscript(i,m));
        Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
            [&] (const int r)
            {
                tmp[r] *= row[r];
            });
    }
}

Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
    [&] (const int r)
{
    Kokkos::atomic_add(&v.entry(k,r), tmp[r]);
});
```

Adding Scratch Space and Vector Parallelism

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typedef Kokkos::TeamPolicy<ExecSpace> Policy;
const int nnz = X.nnz();
const int VectorSize = 8;
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Policy policy(nnz,TeamSize,VectorSize);

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    ExecSpace::scratch_memory_space,
    Kokkos::MemoryUnmanaged > ScratchSpace;
const size_t bytes = ScratchSpace::shmem_size(R);

Kokkos::parallel_for(
    policy.set_scratch_size(0,Kokkos::PerTeam(bytes)),
    KOKKOS_LAMBDA(Policy::member_type team)
{
    const size_t i = team.league_rank();

    ScratchSpace tmp(team.team_scratch(0), R);

    const size_t k = X.subscript(i,n);
    const double x_val = X.value(i);
```

```
    Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
        [&] (const int r)
    {
        tmp[r] = x_val;
    });

    for (int m = 0; m < N; ++m) {
        if (m != n) {
            const double *row = u[m].rowptr(X.subscript(i,m));
            Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
                [&] (const int r)
            {
                tmp[r] *= row[r];
            });
        }
    }

    Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
        [&] (const int r)
    {
        Kokkos::atomic_add(&v.entry(k,r), tmp[r]);
    });
});
```

■ Final MTTKRP algorithm:

- Incorporates team size > 1 on GPU for better occupancy (process blocks of non-zeros)
- Blocking on R for better cache performance
 - Template algorithm on R block size for better vectorization
 - Choose block size to balance occupancy, tensor reads, minimize remainder term

Adding Team Parallelism

```
typedef Kokkos::TeamPolicy<ExecSpace> Policy;
const bool is_cuda = std::is_same<ExecSpace,
    Kokkos::Cuda>::value;
const int nnz = X.nnz();
const int VectorSize = 8;
const int TeamSize = is_cuda ? 128/VectorSize : 1;
const int LeagueSize = (nnz+TeamSize-1)/TeamSize;
Policy policy(LeagueSize,TeamSize,VectorSize);

typedef Kokkos::View< double**, Kokkos::LayoutRight,
    ExecSpace::scratch_memory_space,
    Kokkos::MemoryUnmanaged > ScratchSpace;
const size_t bytes = ScratchSpace::shmem_size(TeamSize,R);

Kokkos::parallel_for(
    policy.set_scratch_size(0,Kokkos::PerTeam(bytes)),
    KOKKOS_LAMBDA(Policy::member_type team)
{
    const int team_index = team.team_rank();
    const size_t i = team.league_rank()*TeamSize+team_index;
    if (i >= nnz) return;

    ScratchSpace tmp(team.team_scratch(0), TeamSize, R);

    const size_t k = X.subscript(i,n);
    const double x_val = X.value(i);
```

```
    Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
        [&] (const int r)
    {
        tmp(team_index,r) = x_val;
    });

    for (int m = 0; m < n; ++m) {
        if (m != n) {
            const double *row = u[m].rowptr(X.subscript(i,m));
            Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
                [&] (const int r)
            {
                tmp(team_index,r) *= row[r];
            });
        }
    }

    Kokkos::parallel_for(Kokkos::ThreadVectorRange(team,R),
        [&] (const int r)
    {
        Kokkos::atomic_add(&v.entry(k,r), tmp(team_index,r));
    });
});
```