

EFFICIENT ALGORITHMS FOR CONTINGENCY ANALYSIS OF POWER GRIDS

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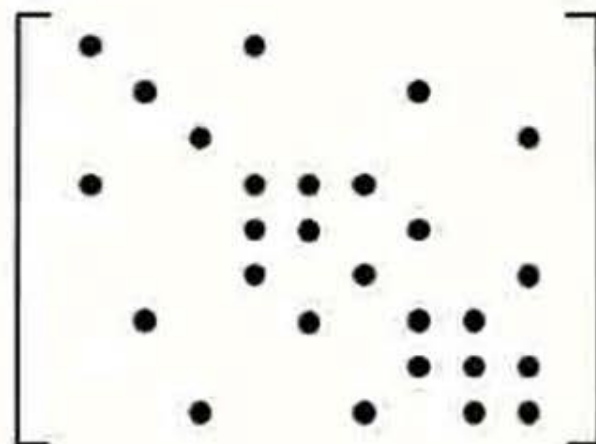
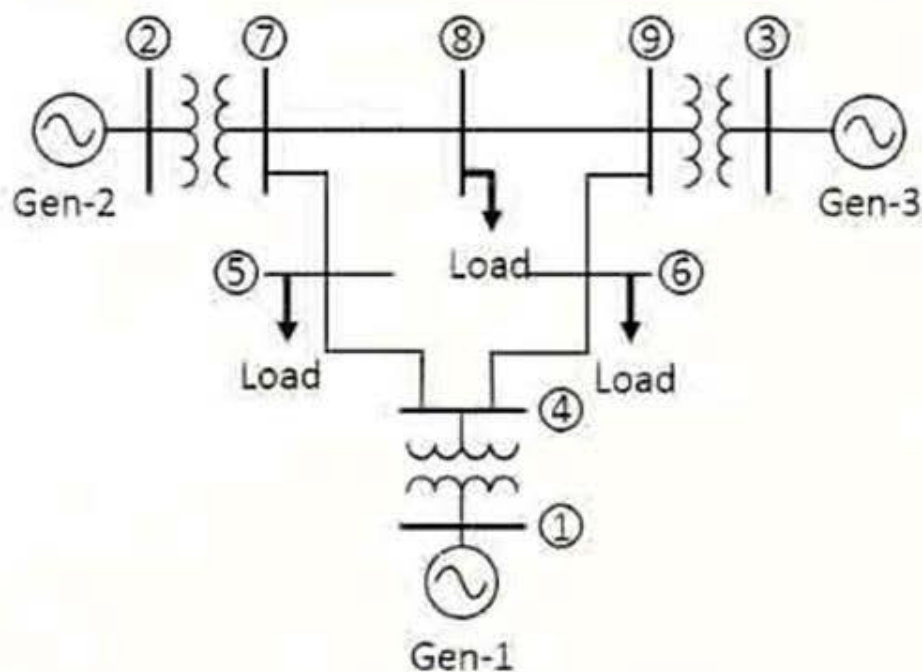
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CONTINGENCY ANALYSIS

CONTINGENCY ANALYSIS

- $N - x$ contingency analysis on power flow evaluates the stability of a power system by simulating the failures of x transmission lines or generators. (N is the number of buses and x is the number of removed components)



CONTINGENCY ANALYSIS

- Currently each $N - x$ contingency analysis is performed independently, but since the number of cases increases exponentially with x , this is computationally impractical. (For example, if $N = 3120$ and $x = 5$, there are in total 2.46×10^{15} cases.)
- Heuristics are needed for choosing the cases to be analyzed. To be able to analyze more cases, we have to speed up the solution time for each case.

CONTINGENCY ANALYSIS

- We propose new algorithms for the problem by observing that only a small portion of the system is changed when a component is removed.
- We first analyze using the “DC” approximate equation:

$$B' \Delta\theta = P$$

where

B' is the imaginary part of the admittance matrix,
 $\Delta\theta$ is the angular separation across a transmission circuit,
 P is the real power flow.

AUGMENTED FORMULATION

AUGMENTED FORMULATION

Inspired by a surgery simulation project:

- Cutting a graphical mesh model of an organ to simulate a surgery using underlying physics-based system of equations.
- The matrix of the system changes while the mesh is being cut.
- Need fast update for interactive simulation.

Reference:

- [Y.-H. Yeung, J. Crouch and A. Pothan](#), Interactively Cutting and Constraining Vertices in Meshes using Augmented Matrices. (Accepted by ACM Transaction on Graphics, under revision)

AUGMENTED FORMULATION

General ideas:

- If only few columns of the matrix change over time, preserve the original matrix and its factors while the changes occur
- Form equivalent **augmented system** of equations with appended rows and columns to compute the solutions of the changed system
- Use a **hybrid (direct + implicit Krylov)** solver or a **direct** solver for the augmented system of equations
- Exploit **sparsity** carefully in the implicit Krylov solver

SOLUTION METHOD

The augmented system of equations can then be expressed as

$$\begin{bmatrix} K & J \\ H & 0 \end{bmatrix} \begin{bmatrix} \alpha_1 \\ \alpha_2 \end{bmatrix} = \begin{bmatrix} f_G \\ 0 \end{bmatrix}$$

- α_1 is the solution vector of the unaffected buses.
- α_2 is the solution vector of the buses affected by the removal of components.

EXPLOITING SPARSITY IN VECTORS

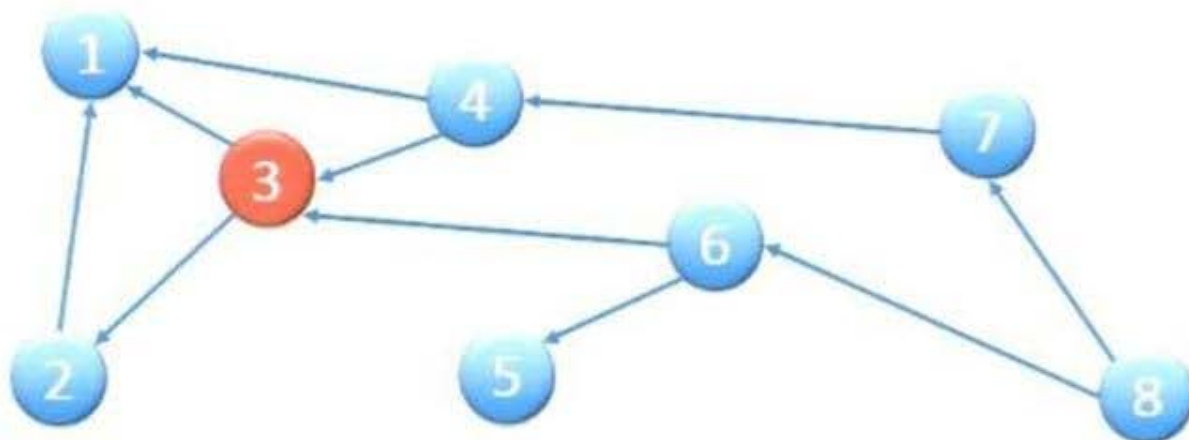
EXPLOITING SPARSITY IN VECTORS

Consider the equation $z = Hy = H(LDL^T)^{-1}b$:

- Sparsity in b in solving with L (closure of nonzero indices of f in $G(L)$, will be defined in the next slide)
- H is a submatrix of an identity matrix, so $z = Hy$ selects few elements of y . Need to compute only components of y selected by H in solving with L^T (closure of selected components of y in $G(L)$)

EXPLOITING SPARSITY IN VECTORS

Closure in of a vector $v = e_3$ in a graph $G(L)$:

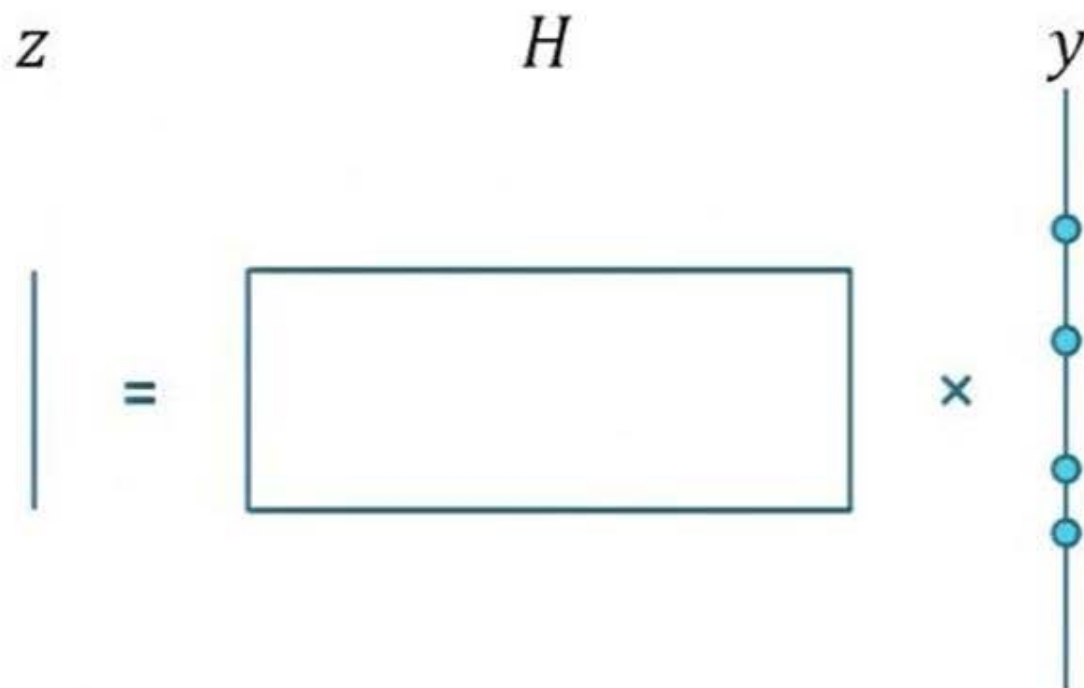


Reference:

- J. Gilbert, "Predicting structure in sparse matrix computations," SIAM Journal on Matrix Analysis and Applications, vol. 15, pp. 62–79, 1994.

SPARSITY IN THE MATRIX H

Finding the needed entries of the vector y which are included in the vector z

$$z = H y$$


SPARSITY: NEEDED SOLUTION COMPONENTS

Partial backward substitution by only computing columns of L^T induced by the needed entries of the solution vector



EXPLOITING SPARSITY

- For the GMRES step, copy submatrices of L and L^T for use in multiple iterations (for memory performance)
- For other steps, compute indices of nonzero components of the vectors to identify the rows and columns of L and L^T needed

USING DIRECT METHOD

When the augmented size is small, we could form the matrix $H(LDL^T)^{-1}J$ explicitly as follows:

$$\text{Let } S = H(LDL^T)^{-1}J = HL^{-T}D^{-1}L^{-1}J.$$

1. Solve $LZ = J$ for Z .
2. Solve $L^TY = H$ for Y .
3. Compute $S = Y^TD^{-1}Z$.

Sparsity can be exploited in solving Steps 1 & 2. All 3 steps can be implemented in parallel.

TIME COMPLEXITY

- Initialization step:
 1. Compute LDL^T factorization of K_G ($O(n^3)$ for general matrix or $O(n^{3/2})$ for planar graph)
- Update step:
 1. Update K_G . ($O(m)$)
 2. Compute matrices J and H . ($O(m)$)
 3. Solve $Ky = f_G$ and compute $z = Hy$. ($\sigma = O(|\text{closure}(f_G)| + |\text{closure}(\hat{y})|) \leq O(\text{nnz}(L))$ where \hat{y} is the components of y selected by H)

TIME COMPLEXITY

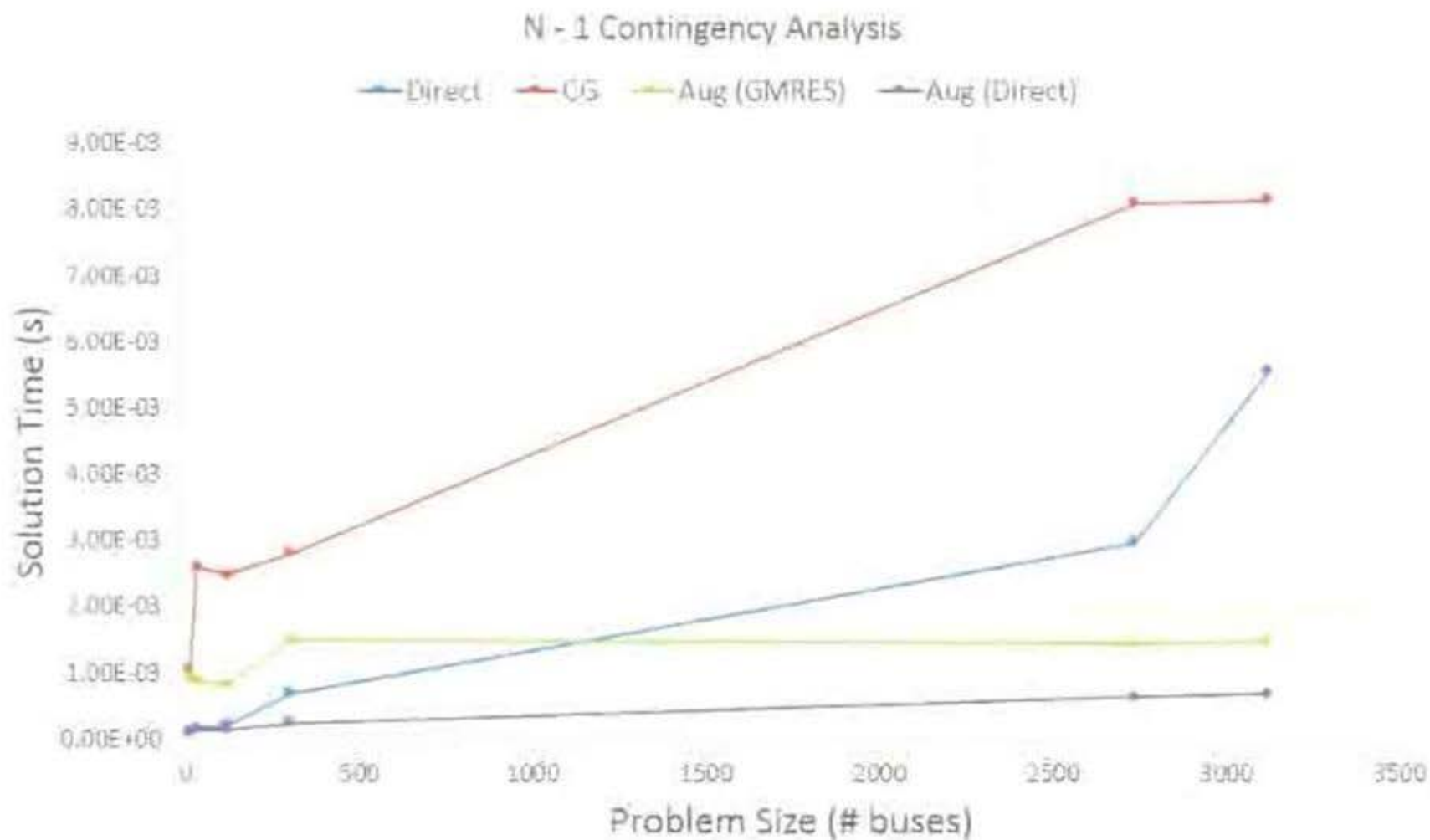
4. Solve $(HK^{-1}J)\alpha_2 = z$. ($O(\sigma \cdot n_{iter})$ for implicit GMRES where n_{iter} is the number of iterations; or $O(\sigma \cdot m + m^2 \cdot n + m^2)$ for serial implementation, $O(\sigma + n + m^2)$ for parallel implementation of direct method.
 5. Solve $K\alpha_1 = f_G - J\alpha_2$. ($O(\text{nnz}(L))$)
- Upper bound complexity is $O(\text{nnz}(L) \cdot n_{iter})$ for using GMRES implicitly.

RESULTS

TEST CASES

Problem size	nnz(B')	Condition number
9	24	8.99×10^1
30	107	9.62×10^2
118	463	4.85×10^3
300	1115	1.55×10^5
2736	9251	1.09×10^6
3120	10477	1.25×10^6

N - 1 CONTINGENCY ANALYSIS



EXTENSION TO AC POWER FLOW

EXTENSION TO AC POWER FLOW

Extension to AC Power Flow:

$$Y\mathbf{v} = \mathbf{b}^*(\mathbf{v})$$
$$\mathbf{b}^*(\mathbf{v})^T = \begin{bmatrix} \frac{s_1^*}{v_1^*} & \frac{s_2^*}{v_2^*} & \dots & \frac{s_n^*}{v_n^*} \end{bmatrix}$$

Where

Y is the admittance matrix,
 v_i is the voltage at each bus i ,
 s_i is the power at each bus i .

All values are complex.

EXTENSION TO AC POWER FLOW

Let $F(\mathbf{v}) = Y\mathbf{v} - \mathbf{b}^*(\mathbf{v}) = 0$.

Newton's method:

$$F'_k \Delta \mathbf{v} = -F(\mathbf{v}_k)$$

where

$$\Delta \mathbf{v} = \mathbf{v}_{k+1} - \mathbf{v}_k.$$

EXTENSION TO AC POWER FLOW

Equation becomes

$$[j(Y + S^*) \quad Y + S^*] \Delta \mathbf{y} = -F(\mathbf{v}_k)$$

where

$$S = \text{diag}(\mathbf{s}) \text{diag}(\mathbf{v}_k)^{-2},$$
$$\Delta \mathbf{y} = \begin{bmatrix} \text{diag}(\mathbf{v}_k) \\ \text{diag}(\mathbf{v}_k) \text{diag}(\boldsymbol{\xi}_k)^{-1} \end{bmatrix} \Delta \mathbf{x}.$$

EXTENSION TO AC POWER FLOW

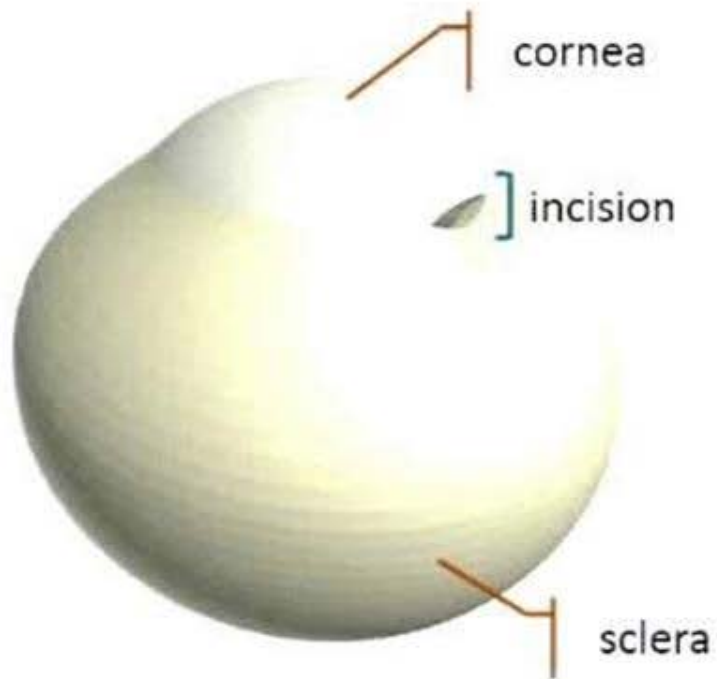
Separating real and imaginary parts,

$$\underbrace{\begin{bmatrix} -\Im(Y) & \Re(Y) \\ \Re(Y) & \Im(Y) \end{bmatrix}}_{\text{Only in } Y; \text{ constant across iterations}} \Delta \mathbf{y}_{k+1} = \underbrace{\begin{bmatrix} -\Im(S) & -\Re(S) \\ -\Re(S) & \Im(S) \end{bmatrix}}_{\text{changes every iteration; involves only matrix-vector multiplication}} \Delta \mathbf{y}_k - \begin{bmatrix} \Re(F(\mathbf{v}_k)) \\ \Im(F(\mathbf{v}_k)) \end{bmatrix}$$

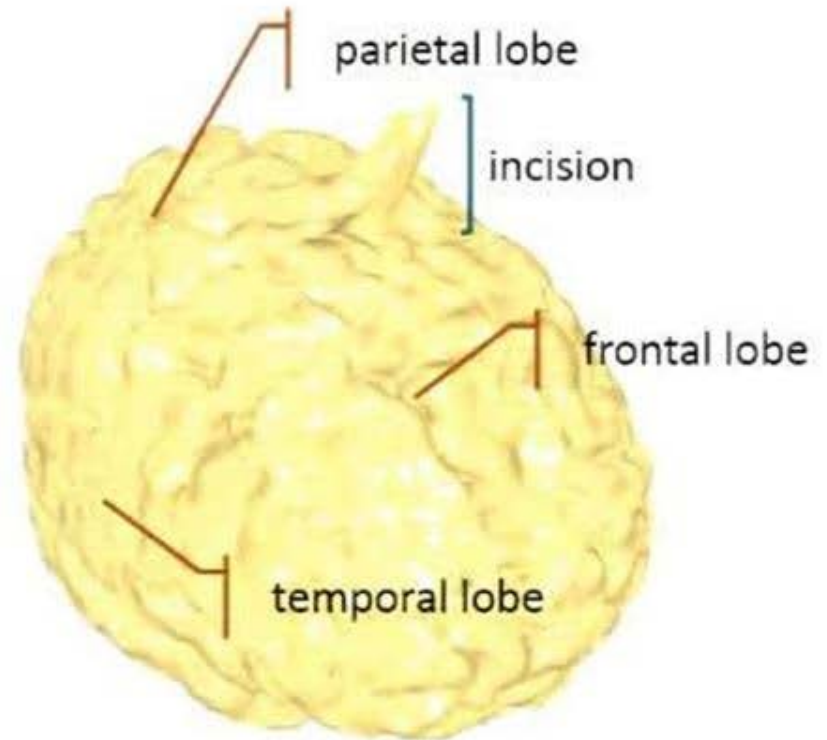
- We are currently working on the convergence problem.

RESULTS OF SURGERY SIMULATION

SURGERY SIMULATION



Eye Model

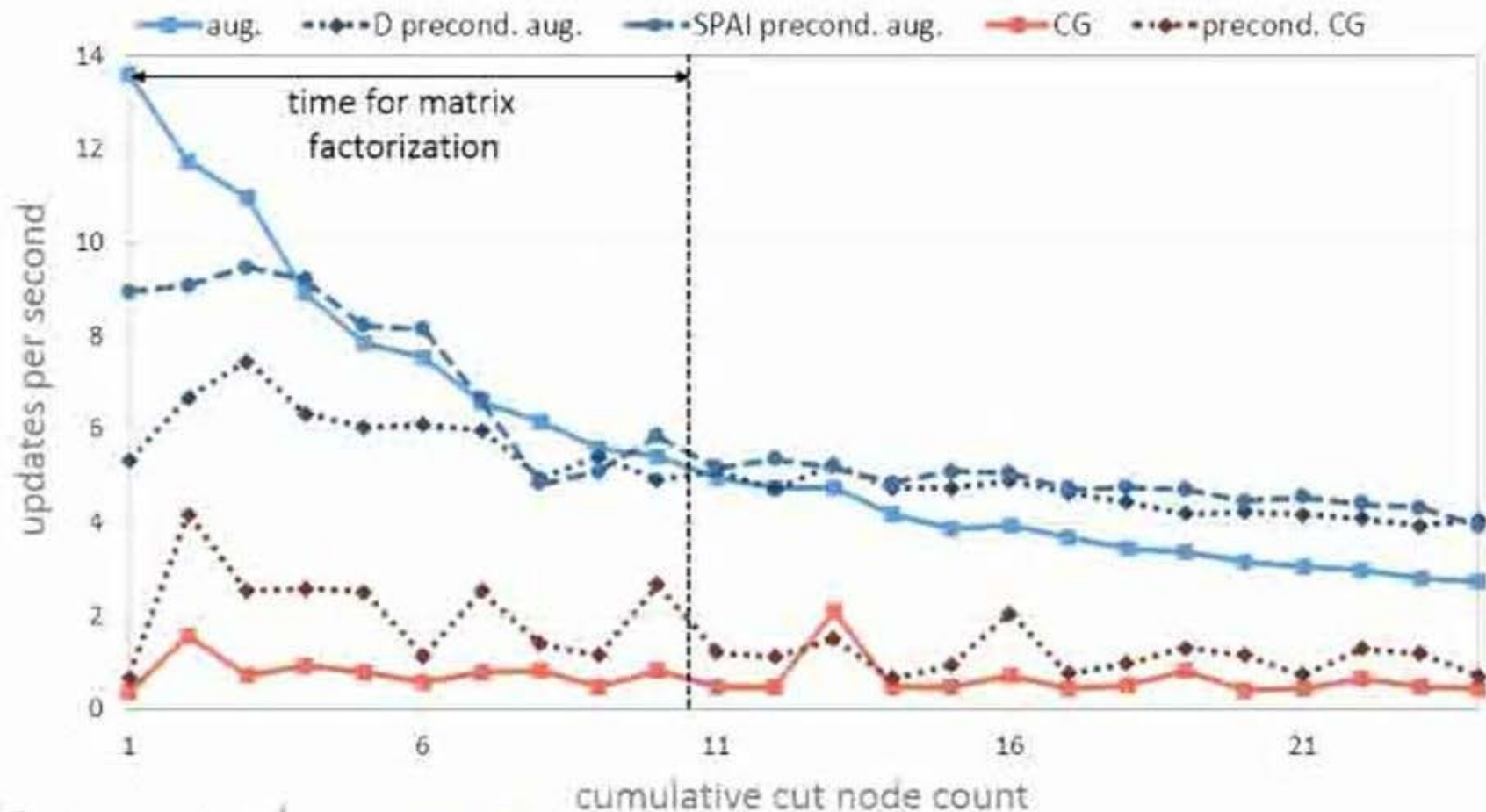


Brain Model

SURGERY SIMULATION

Condition number: 1.62×10^7

Cutting of Eye Mesh: 49,674 elements and 16,176 Nodes



SUMMARY

- We propose an augmented formulation for solving systems of equations undergoing small changes with time complexity linear in the number of nonzeros of the factors.
- Results from surgery simulation give promising speedups over conjugate gradient (CG) method.
- The admittance matrix across cases in contingency analysis of power flow system is quite similar to the stiffness matrix in finite element model.
- We showed that we can solve the problem about 8 times faster than direct and CG methods. Hence we can solve many more contingency analyses in a given time.

QUESTIONS & ANSWERS

THANK YOU!