# Beyond the Black Box in Derivative-Free and Simulation-Based Optimization 

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## Optimizing (Almost) Everything!

0. Simulation-based and derivative-free optimization?
I. Optimization of black boxes
$\diamond$ Empirical performance tuning of HPC codes
$\diamond$ Model-based algorithms
II. Exploiting structure in functions of black boxes
$\diamond$ Least squares - calibrating DFT sims
$\diamond$ Nonsmoothness - bioremediation
$\diamond$ Some partials - multilevel energy functionals
$\diamond$ Constraints

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## Doing <br> something <br> with <br> little

Doing
better
with
little
more


## Simulation-Based Optimization

$$
\min _{x \in \mathbb{R}^{n}}\left\{f(x)=F[x, S(x)]: c_{I}[x, S(x)] \leq 0, c_{E}[x, S(x)]=0\right\}
$$

$\approx$ "parameter estimation" $\approx$ "model calibration" $\approx$ "design optimization" $\approx \ldots$
$\diamond S: \mathbb{R}^{n} \rightarrow \mathbb{C}^{p}$ simulation output, often "noisy" (even when deterministic)
$\diamond$ Derivatives $\nabla_{x} S$ often unavailable or
prohibitively expensive to obtain/approximate directly
$\diamond S$ can contribute to objective and/or constraints
$\diamond$ Single evaluation of $S$ could take seconds/minutes/hours/...
$\Rightarrow$ Evaluation is a bottleneck for optimization


## Blame Computing!

## for pervasiveness of simulations in sci\&eng

$\diamond$ Parallel/multi-core environments common
$\diamond$ Simulations ("forward problem") faster, more realistic/complex


Argonne's AVIDAC
(1953 vacuum tubes)


Argonne's BlueGene/Q (2012 0.8M cores)


Sunway TaihuLight (2016 11M cores)

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## for the challenges in SBO

$\diamond$ Optimization, UQ often an afterthought
$\diamond$ Obstacles for Algorithmic Differentiation (coupled legacy/proprietary codes, memory)
$\rightarrow$ [Coleman \& Xu; SIAM 2016], [Griewank \& Walther; SIAM 2008] $\rightarrow$ MS76 today!
$\diamond$ Computational noise can complicate everything

$$
\rightarrow \text { [Moré \& W.; SISC 2011] }
$$

$\diamond$ Finite differences noisy, possibly expensive
$\rightarrow$ [Moré \& W.; TOMS 2012]
$\diamond$ Computational budget limits \# f evals

## Derivative-Free Optimization

"Some derivatives $\left(\nabla_{x} S(x)\right)$ unavailable for optimization purposes"

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## The Challenge: <br> Optimization is tightly coupled with derivatives

Typical optimality (no noise, smooth functions)

$$
\nabla_{x} f\left(x_{*}\right)+\lambda^{\top} \nabla_{x} c_{E}\left(x_{*}\right)=0, c_{E}\left(x_{*}\right)=0
$$



William Karush [Optimization Stories, 2012]

(sub)gradients $\nabla_{x} f, \nabla_{x} c$ enable:
$\diamond$ Faster feasibility
$\diamond$ Faster convergence

- Guaranteed descent
- Approximation of nonlinearities
$\diamond$ Better termination
- Measure of criticality $\left\|\nabla_{x} f\right\|,\left\|\mathcal{P}_{\Omega}\left(\nabla_{x} f\right)\right\|$
$\diamond$ Sensitivity analysis
- Correlations, standard errors, UQ, ...


## The Price of Algorithm Choice: Solvers in PETSc/TAO



Toolkit for Advanced Optimization
[Munson et al.; mcs.anl.gov/tao]

## Increasing level of user input:

$n m$ Assumes $\nabla_{x} f$ unavailable, black box
pounders Assumes $\nabla_{x} f$ unavailable, exploits problem structure

Imvm Uses available $\nabla_{x} f$

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THIS TALK
Imvm Uses available $\nabla_{x} f$

DFO methods should be designed to beat finite-difference-based methods

Observe: Constrained by budget on \#evals, method limits solution accuracy/problem size


## Black-Box Optimization



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## Black-Box Optimization

Inputs

$$
\min _{x \in \mathbb{R}^{n}} f(x)
$$

## Only access to $f=S$ is through sampling

$\diamond$ (Scalar) Output of an experiment
$\diamond$ Proprietary libraries/closed codes
$\diamond$ Often discrete/compact domains
Throughout this talk:
"Black box" is both good and evil

## A Black Box: Automating Empirical Performance Tuning

Given semantically equivalent codes $x_{1}, x_{2}, \ldots$, minimize run time subject to energy consumption



## $\min \left\{f(x):\left(x_{\mathcal{C}}, x_{\mathcal{I}}, x_{\mathcal{B}}\right) \in \Omega_{\mathcal{C}} \times \Omega_{\mathcal{I}} \times \Omega_{\mathcal{B}}\right\}$

$x$ multidimensional parameterization (compiler type, compiler flags, unroll/tiling factors, internal tolerances, ...)
$\Omega$ search domain (feasible transformation, no errors)
$f$ quantifiable performance objective (requires a run)
$\rightarrow$ [Audet \& Orban; SIOPT 2006], [Balaprakash, W., Hovland; ICCS 2011], [Porcelli \& Toint; 2016] Numerical Linear Algebra $\rightarrow$ [N. Higham; SIMAX 1993], . . .

## Black-Box Algorithms: Stochastic Methods

## Random search

## Repeat:

1. Randomly generate direction $d_{k} \in \mathbb{R}^{n}$
2. Evaluate "gradient-free oracle" $g\left(x_{k} ; h_{k}\right)=\frac{f\left(x_{k}+h_{k} d_{k}\right)-f\left(x_{k}\right)}{h_{k}} d_{k}$

$$
\text { ( } \approx \text { directional derivative) }
$$

3. Compute $x_{k+1}=x_{k}-\delta_{k} g\left(x_{k} ; h_{k}\right)$, evaluate $f\left(x_{k+1}\right)$

Convergence (for different types of $f$ ) tends to be probabilistic [Kiefer \& Wolfowitz; AnnMS 1952], [Polyak; 1987], [Ghadimi \& Lan; SIOPT 2013], [Nesterov \& Spokoiny; FoCM 2015], . . .

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## Stochastic heuristics (nature-inspired methods, etc.)

$\diamond$ Popular in practice, especially in engineering
$\diamond$ Typically global in nature
$\diamond$ Require many $f$ evaluations

## Black-Box Algorithms: Direct Search Methods

Pattern Search + Variants


Easy to parallelize $f$ evaluations

## Nelder-Mead + Variants



Popularized by Numerical Recipes
$\diamond$ Rely on indicator functions: $\left[f\left(x_{k}+s\right)<? f\left(x_{k}\right)\right]$
$\diamond$ Work with black-box $f(x)$, do not exploit structure $F[x, S(x)]$
$\diamond$ Convergence results for variety of settings

Survey $\rightarrow$ [Kolda, Lewis, Torczon; SIREV 2003]
Newer NM $\rightarrow$ [Lagarias, Poonen, Wright; SIOPT 2012] Tools $\rightarrow$ DFL [Liuzzi et al.], NOMAD [Audet et al.], . . .

## Making the Most of Little Information About Smooth $f$

$\diamond$ Overhead of the optimization routine is minimal (negligible?) relative to cost of evaluating simulation


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## Bank of data, $\left\{x_{i}, f\left(x_{i}\right)\right\}_{i=1}^{k}$ :

$=$ Points (\& function values) evaluated so far
$=$ Everything known about $f$
Goal:
$\diamond$ Make use of growing Bank as optimization progresses
$\diamond$ Limit unnecessary evaluations
(geometry/approximation)

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## Model-Based Trust-Region Algorithms

Substitute $\min \left\{q_{k}(x): x \in \mathcal{B}_{k}\right\}$ for $\min f(x)$

$$
q_{k}(x)=f\left(x_{k}\right)+g_{k}^{\top}\left(x-x_{k}\right)+\frac{1}{2}\left(x-x_{k}\right)^{\top} H_{k}\left(x-x_{k}\right)
$$

$f$ expensive, no $\nabla f$
$q_{k}$ cheap, analytic derivatives


## Trust region:

$\mathcal{B}_{k}=\left\{x \in \Omega:\left\|x-x_{k}\right\| \leq \Delta_{k}\right\}$
$\diamond$ Trust $q_{k} \approx f$ in $\mathcal{B}_{k}$
$\diamond$ Update based on $\rho_{k}=\frac{f\left(x_{k}\right)-f\left(x_{+}\right)}{q_{k}\left(x_{k}\right)-q_{k}\left(x_{+}\right)}$

## Typical models

$\diamond$ Only need (occasional) local approximation
$\diamond$ Taylor-based: $g_{k}=\nabla f\left(x_{k}\right)$, $H_{k} \approx \nabla^{2} f\left(x_{k}\right)$
$\rightarrow$ [Conn, Gould, Toint; SIAM 2000]

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## Black-Box Algorithms: Building Models Without Derivatives

Given data $\left(\mathcal{X}_{k}, f\left(\mathcal{X}_{k}\right)\right)$ and basis $\Phi$, "solve"

$$
\Phi\left(\mathcal{X}_{k}\right) z=\left[\begin{array}{lll}
\Phi_{c} & \Phi_{g} & \Phi_{H}
\end{array}\right]\left[\begin{array}{c}
z_{c} \\
z_{g} \\
z_{H}
\end{array}\right]=\underline{\mathrm{f}}=f\left(\mathcal{X}_{k}\right)
$$

## Full quadratics, $\left|\mathcal{X}_{k}\right|=\frac{(n+1)(n+2)}{2}$

$\diamond$ Interpolation: $q_{k}\left(y_{i}\right)=f\left(y_{i}\right), \quad \forall y_{i} \in \mathcal{X}_{k}$
$\diamond$ Geometric conditions on points in $\mathcal{X}_{k}$

## Undetermined interp., $\left|\mathcal{X}_{k}\right|<\frac{(n+1)(n+2)}{2}$

$\diamond$ Use (Powell) Hessian updates

$$
\begin{aligned}
\min _{g_{k}, H_{k}} & \left\|H_{k}-H_{k-1}\right\|_{F}^{2} \\
\text { s.t. } & q_{k}=\underline{\mathrm{f}} \text { on } \mathcal{X}_{k}
\end{aligned}
$$

$$
\text { Regression, }\left|\mathcal{X}_{k}\right|>\frac{(n+1)(n+2)}{2}
$$

$\diamond$ Solve $\min _{z}\|\Phi z-\underline{f}\|$

## Multivariate (Scattered Data) Interpolation is a Different Kind of Animal

$$
m\left(y_{i}\right)=f\left(y_{i}\right) \quad \forall y_{i} \in \mathcal{X}
$$

$n=1$ Given distinct points, can find a unique degree $|\mathcal{X}|-1$ polynomial $m$
$n>1$ Not true! (see Mairhuber-Curtis Theorem)

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6th point for a quadratic in $\mathbb{R}^{2}$


Nearby constraints affect geometry
$\rightarrow$ [Wendland; CUP 2010]

## Convergence to Stationary Points \& Software

## $\lim _{k \rightarrow \infty} \nabla f\left(x_{k}\right)=0$ provided:

$0 . f$ is sufficiently smooth and regular (e.g., bounded level sets)

1. Control $\mathcal{B}_{k}$ based on model quality
2. (Occasional) approximation within $\mathcal{B}_{k}$

Our quadratics satisfy

- $\left|q_{k}(x)-f(x)\right| \leq \kappa_{1}\left(\gamma_{f}+\left\|H_{k}\right\|\right) \Delta_{k}^{2}, \quad \forall x \in \mathcal{B}_{k}$
- $\left\|g_{k}+H_{k}\left(x-x_{k}\right)-\nabla f(x)\right\| \leq \kappa_{2}\left(\gamma_{f}+\left\|H_{k}\right\|\right) \Delta_{k}, \quad \forall x \in \mathcal{B}_{k}$

3. Sufficient decrease


Survey $\rightarrow$ [Conn, Scheinberg, Vicente; SIAM 2009] Methods $\rightarrow$ [Powell: COBYLA, UOBYQA, NEWUOA, BOBYQA, LINCOA],

Line search methods also work $\rightarrow$ [Kelley et al; IFFCO] RBF models also work $\rightarrow$ [W. \& Shoemaker; SIREV 2013] Probabilistic models $\rightarrow$ [Bandeira, Scheinberg, Vicente; SIOPT 2014]

Michael J.D. Powell, 1936-2015


## Structure in Simulation-Based Optimization, $\min f(x)=F[x, S(x)]$

$f$ is often not a black box $S$
NLS Nonlinear least squares

$$
f(x)=\sum_{i}\left(S_{i}(x)-d_{i}\right)^{2}
$$

CNO Composite (nonsmooth) optimization

$$
f(x)=h(S(x))
$$

SKP Not all variables enter simulation

$$
f(x)=g\left(x_{I}, x_{J}\right)+h\left(S\left(x_{J}\right)\right)
$$

SCO Only some constraints depend on simulation

$$
\min \left\{f(x): c_{1}(x)=0, c_{S}(x)=0\right\}
$$

+ Slack variables

$$
\Omega_{S}=\left\{\left(x_{I}, x_{J}\right): S\left(x_{J}\right)+x_{I}=0, x_{I} \geq 0\right\}
$$

Model-based methods offer one way to exploit such structure

## General Setting - Modeling Smooth $S_{1}(x), S_{2}(x), \ldots, S_{p}(x)$

## Assume:

$\bigcirc$ each $S_{i}$ is continuously differentiable, available
$\diamond$ each $\nabla S_{i}$ is Lipschitz continuous, unavailable

## General Setting - Modeling Smooth $S_{1}(x), S_{2}(x), \ldots, S_{p}(x)$

## Assume:

$\diamond$ each $S_{i}$ is continuously differentiable, available
$\diamond$ each $\nabla S_{i}$ is Lipschitz continuous, unavailable
$m^{S_{i}}: \mathbb{R}^{n} \rightarrow \mathbb{R}$ approximates $S_{i}$ on $\mathcal{B}(x, \Delta)$ $i=1, \ldots, p$

## Fully Linear Models

$m^{S_{i}}$ fully linear on $\mathcal{B}(x, \Delta)$ if there exist constants $\kappa_{i, \text { ef }}$ and $\kappa_{i, \text { eg }}$ independent of $x$ and $\Delta$ so that

$$
\begin{aligned}
\left|S_{i}(x+s)-m^{S_{i}}(x+s)\right| \leq \kappa_{i, \text { ef }} \Delta^{2} & \forall s \in \mathcal{B}(0, \Delta) \\
\left\|\nabla S_{i}(x+s)-\nabla m^{S_{i}}(x+s)\right\| \leq \kappa_{i, \mathrm{eg}} \Delta & \forall s \in \mathcal{B}(0, \Delta)
\end{aligned}
$$

## NLS- Nonlinear Least Squares $f(x)=\frac{1}{2} \sum_{i} R_{i}(x)^{2}$

## Obtain a vector of output $R_{1}(x), \ldots, R_{p}(x)$

$\diamond$ Model each $R_{i}$

$$
R_{i}(x) \approx m_{k}^{R_{i}}(x)=R_{i}\left(x_{k}\right)+\left(x-x_{k}\right)^{\top} g_{k}^{(i)}+\frac{1}{2}\left(x-x_{k}\right)^{\top} H_{k}^{(i)}\left(x-x_{k}\right)
$$

$\diamond$ Approximate:

$$
\begin{aligned}
\nabla f(x)= & \sum_{i} \nabla \mathbf{R}_{\mathbf{i}}(\mathbf{x}) R_{i}(x) \quad \longrightarrow \sum_{i} \nabla m_{k}^{R_{i}}(x) R_{i}(x) \\
\nabla^{2} f(x)= & \sum_{i} \nabla \mathbf{R}_{\mathbf{i}}(\mathbf{x}) \nabla \mathbf{R}_{\mathbf{i}}(\mathbf{x})^{\top}+\sum_{i} R_{i}(x) \nabla^{\mathbf{2}} \mathbf{R}_{\mathbf{i}}(\mathbf{x}) \\
& \longrightarrow \sum_{i} \nabla m_{k}^{R_{i}}(x) \nabla m_{k}^{R_{i}}(x)^{\top}+\sum_{i} R_{i}(x) \nabla^{2} m_{k}^{R_{i}}(x)
\end{aligned}
$$

$\diamond$ Model $f$ via Gauss-Newton or similar

> regularized Hessians $\rightarrow$ DFLS [Zhang, Conn, Scheinberg] full Newton $\rightarrow$ POUNDERS [W., Moré]

## NLS- Consequences for $f(x)=\frac{1}{2} \sum_{i} R_{i}(x)^{2}$

Pay a (negligible for expensive $S$ ) price in terms of $p$ models
$\diamond$ Save linear algebra using interpolation set $\mathcal{X}_{k}$ common to all models

- Single system solve, multiple right hand sides

$$
\Phi\left(\mathcal{X}_{k}\right)\left[\begin{array}{lll}
z^{(1)} & \cdots & z^{(p)}
\end{array}\right]=\left[\begin{array}{lll}
\mathrm{R}_{1} & \cdots & \underline{\mathrm{R}}_{p}
\end{array}\right]
$$

- $m^{R_{1}}$ quality $\Rightarrow$ quality of all $m^{R_{i}}$
+ (nearly) exact gradients for $R_{i}$ (nearly) linear
- No longer interpolate function at data points

$$
\begin{aligned}
m\left(x_{k}+\delta\right)= & f\left(x_{k}\right) \\
& +\delta^{\top} \sum_{i} g_{k}^{(i)} R_{i}\left(x_{k}\right) \\
& +\frac{1}{2} \delta^{\top} \sum_{i}\left(g_{k}^{(i)}\left(g_{k}^{(i)}\right)^{\top}+R_{i}\left(x_{k}\right) H_{k}^{(i)}\right) \delta \\
& + \text { missing h.o. terms }
\end{aligned}
$$

## NLS- POUNDERS in Practice: DFT Calibration/MLE

$\min _{x} \sum_{i=1}^{p} w_{i}\left(S_{i}(x)-d_{i}\right)^{2}$
$S_{i}(x)$ Simulated (DFT) nucleus property
$d_{i}$ Experimental data $i$
$w_{i}$ Weight for data type $i$
$p$ Parallel simulations ( 12 wallclock mins)

$\rightarrow$ [Kortelainen et al., PhysRevC 2010]


Energy Residual [MeV], Nucleus \#22


## CNO- Composite Nonsmooth Optimization Examples

## Ex.- Groundwater remediation

Determine rates $x$ for extraction/injection wells
$\diamond$ Regulator's simulator returns flow $S_{i}(x)$ in/out of cell $i$
$\diamond$ Minimize plume fluxes (e.g., regulatory \$ penalties) $f(x)=\sum_{i}\left|S_{i}(x)\right|$
$\rightarrow$ See MS90 later today


## Ex.- Particle accelerator design

Minimize particle losses: $f(x)=\max _{t_{i} \in \mathcal{T}_{1}} S\left(x ; t_{i}\right)-\min _{t_{i} \in \mathcal{T}_{2}(x)} S\left(x ; t_{i}\right)$

CNO- Composite Nonsmooth Optimization $f(x)=h(S(x) ; x)$
nonsmooth (algebraically available) function $h: \mathbb{R}^{p} \times \mathbb{R}^{n} \rightarrow \mathbb{R}$ of a smooth (blackbox) mapping $S: \mathbb{R}^{n} \rightarrow \mathbb{R}^{p}$

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nonsmooth (algebraically available) function $h: \mathbb{R}^{p} \times \mathbb{R}^{n} \rightarrow \mathbb{R}$ of a smooth (blackbox) mapping $S: \mathbb{R}^{n} \rightarrow \mathbb{R}^{p}$

Basic Idea: Knowledge of vector $S\left(x^{k}\right)$ \& potential nondifferentiability at $S\left(x^{k}\right)$ should enhance (theoretical and practical) progress to a stationary point

EX.- $f^{1}(x)=\|S(x)\|_{1}=\sum_{i=1}^{p}\left|S_{i}(x)\right|$

$$
\partial f^{1}(x)=\sum_{i: S_{i}(x) \neq 0} \operatorname{sgn}\left(S_{i}(x)\right) \nabla S_{i}(x)+\sum_{i: S_{i}(x)=0} \operatorname{co}\left\{-\nabla S_{i}(x), \nabla S_{i}(x)\right\}
$$

$\diamond \mathcal{D}^{c}=\left\{x: \exists i\right.$ with $\left.S_{i}(x)=0, \nabla S_{i}(x) \neq 0\right\}$

+ Compact $\partial f(x)$
- $\mathcal{D}^{c}$ depends on $\nabla S_{i}(x)$



## CNO- The Nuisance Set, $\mathcal{N}$

Relaxation $\mathcal{N} \subseteq \mathcal{D}^{c}$ using only zero-order information
$f^{1}$ :

$$
\mathcal{N}=\left\{x: \exists i \text { with } S_{i}(x)=0\right\}
$$

$f^{\infty}$ :

$$
\mathcal{N}=\left\{x: f^{\infty}(x)=0 \text { or }\left|\arg \max _{i}\right| S_{i}(x)| |>1\right\}
$$



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$$



## Observe

When $x^{k} \notin \mathcal{N}$,

$$
\begin{aligned}
\partial f\left(x^{k}\right) & =\nabla_{f}\left(x^{k}\right) \\
& =\nabla_{x} S\left(x^{k}\right)^{\top} \nabla_{S} h\left(S\left(x^{k}\right)\right) \\
& \approx \nabla_{x} M\left(x^{k}\right)^{\top} \nabla_{S} h\left(S\left(x^{k}\right)\right)
\end{aligned}
$$

and smooth approximation is justified


## CNO- Subdifferential Approximation

$\diamond x^{k} \in \mathcal{N}$, we build a set of generators $\mathcal{G}\left(x^{k}\right)$ based on $\partial_{S} h\left(S\left(x^{k}\right)\right)$.

- co $\left\{\mathcal{G}\left(x^{k}\right)\right\}$ approximates $\partial f\left(x^{k}\right)$

$$
\begin{aligned}
& \text { Ex.- } f^{1}(x)=\|S(x)\|_{1} \\
& \mathcal{G}\left(x^{k}\right)=\nabla M\left(x^{k}\right)^{\top}\left\{\operatorname{sgn}\left(S\left(x^{k}\right)\right)+\underset{i: S_{i}\left(x^{k}\right)=0}{\cup}\left\{-e_{i}, 0, e_{i}\right\}\right\}
\end{aligned}
$$

## CNO- Subdifferential Approximation

$\diamond x^{k} \in \mathcal{N}$, we build a set of generators $\mathcal{G}\left(x^{k}\right)$ based on $\partial_{S} h\left(S\left(x^{k}\right)\right)$.

- co $\left\{\mathcal{G}\left(x^{k}\right)\right\}$ approximates $\partial f\left(x^{k}\right)$

$$
\begin{aligned}
& \text { Ex.- } f^{1}(x)=\|S(x)\|_{1} \\
& \mathcal{G}\left(x^{k}\right)=\nabla M\left(x^{k}\right)^{\top}\left\{\operatorname{sgn}\left(S\left(x^{k}\right)\right)+\underset{i: S_{i}\left(x^{k}\right)=0}{\cup}\left\{-e_{i}, 0, e_{i}\right\}\right\}
\end{aligned}
$$

Nearby data $\mathcal{X} \subset \mathcal{B}\left(x^{k}, \Delta_{k}\right)$ informs models $M=m^{S}$ and generator set
$\diamond$ Manifold sampling method uses manifold(s) of $\mathcal{X}$

$$
\nabla M\left(x^{k}\right)^{\top} \underset{y^{i} \in \mathcal{X}}{\cup} \operatorname{mani}\left(S\left(y^{i}\right)\right)
$$

$\diamond$ Traditional gradient sampling
$\rightarrow$ [Burke, Lewis, Overton; SIOPT 2005]

$$
\underset{y^{i} \in \mathcal{X}}{\cup} \nabla M\left(y^{i}\right)^{\top} \operatorname{mani}\left(S\left(y^{i}\right)\right)
$$

## CNO- Smooth Trust-Region Subproblem

Smooth master model from minimum-norm element

$$
m^{f}\left(x^{k}+s\right)=f\left(x^{k}\right)+\left\langle s, \operatorname{proj}\left(0, \operatorname{co}\left\{\mathcal{G}\left(x^{k}\right)\right\}\right)\right\rangle+\cdots
$$

## CNO- Smooth Trust-Region Subproblem

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$$

$\Rightarrow$ smooth subproblems
$\min \left\{m^{f}\left(x^{k}+s\right): s \in \mathcal{B}\left(0, \Delta_{k}\right)\right\}$
$\diamond$ Convex $h$ (e.g., $\|S(x)\|_{1}$ ) and $\nabla S_{i}$ is Lipschitz
$\Rightarrow$ every cluster point of $\left\{x^{k}\right\}_{k}$ is Clarke stationary
$\rightarrow$ [Larson, Menickelly, W.; Preprint 2016]
$\diamond$ OK to sample at $x^{k} \in \mathcal{D}^{C}$
$\diamond$ More general (piecewise differentiable $f$ ) results:
$\rightarrow$ [Larson, Khan, W.; in prog. 2016]
(yesterday in MS155!)

Nonsmooth subproblems
vs.

$$
\min \left\{h\left(M\left(x^{k}+s\right)\right): s \in \mathcal{B}\left(0, \Delta_{k}\right)\right\}
$$

$\diamond$ Requires convex $h$
$\rightarrow$ [Fletcher;
MathProgStudy 1982]
$\rightarrow$ [Grapiglia, Yuan,
Yuan; C\&A Math.
2016]
Complexity results
$\rightarrow$ [Garmanjani,
Júdice, Vicente; SIOPT


Roger Fletcher

## CNO- Example Performance on $L_{1}$ Test Problems

Function Value


Stationary Measure


Smooth black-box methods can fail in practice, even when $\mathcal{D}^{C}$ has measure zero

Numerical tests: $\rightarrow$ [Larson, Menickelly, W.; Preprint 2016]

## SKP- Some Known Partials Example

## Ex.- Bi-level model calibration structure

$$
\min _{x}\left\{f(x)=\sum_{i=1}^{p}\left(S_{i}(x)-d_{i}\right)^{2}\right\}
$$

$S_{i}(x)$ solution to lower-level problem depending only on $x_{J}$

$$
\begin{aligned}
S_{i}(x) & =g_{i}(x)+\min _{y}\left\{h_{i}\left(x_{J} ; y\right): y \in \mathcal{D}_{i}\right\} \\
& =g_{i}(x)+h_{i}\left(x_{J} ; y_{i, *}\left[x_{J}\right]\right)
\end{aligned}
$$

For $x=\left(x_{I}, x_{J}\right)$
$\diamond \nabla_{x_{I}} S_{i}\left(x_{I}, x_{J}\right)$ available
$\diamond \nabla_{x_{J}} S_{i}(x) \approx \nabla_{x_{J}} g_{i}(x)+\nabla_{x_{J}} m^{\tilde{S}_{i}}\left(x_{J}\right)$
$\diamond S_{i}(x)$ continuous and smooth in $x_{I}$
$\diamond g_{i}(x)$ cheap to compute!
$\diamond$ No noise/errors introduced in $g_{i}(x)$

## SKP- Some Known Partials

## $x=\left(x_{I}, x_{J}\right) ;$ have $\frac{\partial f}{\partial x_{I}}$ but not $\frac{\partial f}{\partial x_{J}}$

"Solve"

$$
\Phi z=\underline{\mathrm{f}}
$$

with known $z_{g, I}, z_{H, I}$

$$
\left[\begin{array}{ccc}
\Phi_{c} & \Phi_{g, J} & \Phi_{H, J}
\end{array}\right]\left[\begin{array}{c}
z_{c} \\
z_{g, J} \\
z_{H, J}
\end{array}\right]=\underline{\mathrm{f}}-\Phi_{g, I} z_{g, I}-\Phi_{H, I} z_{H, I}
$$

$\diamond$ Still have interpolation where required
$\diamond$ Effectively lowers dimension to $|J|=n-|I|$ for

- approximation
- model-improving evaluations
- linear algebra
$\diamond \lim _{k \rightarrow \infty} \nabla f\left(x_{k}\right)=0$ as before:
- Guaranteed descent in some directions


## SKP- Numerical Results With Some Partials



Three approaches:

- black box
s exploit least squares
m use $\nabla_{x_{I}}$ derivatives
$\diamond n=16,|I|=3$
$\diamond 5-10$ secs/evaluation

Same algorithmic framework, performance advantages from exploiting structure
$\rightarrow$ [Bertolli, Papenbrock, W., PRC 2012]

## SCO- General Constraints

## $\min \left\{f(x): c_{1}(x)=0, c_{S}(x)=0\right\}$

$\diamond$ Lagrangian (key to optimality conditions):

$$
\begin{aligned}
\nabla L & =\nabla f+\lambda_{1}^{\top} \nabla c_{1}+\lambda_{2}^{\top} \nabla \mathbf{c}_{\mathrm{S}} \\
& \rightarrow \nabla f+\lambda_{1}^{\top} \nabla c_{1}+\lambda_{2}^{\top} \nabla m
\end{aligned}
$$

$\diamond$ Use favorite method: filters, augmented Lagrangian, ...
$\diamond$ Slack variables

- Do not increase effective dimension
- Subproblems can treat separately
- Know derivatives
$\rightarrow$ [Lewis \& Torczon; 2010]
Modified AL methods $\rightarrow$ [Diniz-Ehrhardt, Martínez, Pedroso; C\&A Math. 2011] SBO constraints have unique properties $\rightarrow$ [Le Digabel \& W.; ANL/MCS-P5350-0515 2016]


## SCO- What Constraint Derivatives Buy You

Ex.- Augmented Lagrangian methods, $L_{A}(x, \lambda ; \mu)=f(x)-\lambda^{\top} c(x)+\frac{1}{\mu}\|c(x)\|^{2}$

## $\min _{x}\{f(x): c(x)=0\}$

Four approaches:

1. Penalize constraints
2. Treat $c$ and $f$ both as (separate) black boxes
3. Work with $f$ and $\nabla_{x} c$
4. Have both $\nabla_{x} f$ and $\nabla_{x} c$


$$
n=15,11 \text { constraints }
$$

## OPTIMIZE EVERYTHING

## Mathematically unwrap problems to expose (the deepest) black boxes

$\diamond$ Structure is everywhere, even in legacy-code-driven optimization problems
$\diamond$ Exploiting structure is one way to expand range of optimization to solve grand-challenge problems
$\diamond$ Sacrifice little in convenience

- Output \& model residuals $\left\{r_{i}(x)\right\}_{i}$, not $\|r(x)\|$
- Output \& model constraints $\left\{c_{i}(x)\right\}_{i}$, not a penalty $P(c(x))$
- Explicitly handle nonsmoothness (and noise, ...)
$\diamond$ Papers and links at www.mcs.anl.gov/~wild
$\diamond$ Collaborators in this work:
Awesome opportunities for students
$\begin{array}{ll} & \begin{array}{l}\text { Aswin Kannan (UIUC), Slava Kungurtsev (UCSD), } \\ \text { Jeff Larson (UC-Denver), Matt Menickelly (Lehigh) }\end{array} \\ \text { Argonne } & \begin{array}{l}\text { _.and postdocs! } \\ \text { Prasanna Balaprakash (UL Bruxelles), Kamil Khan (MIT) }\end{array}\end{array}$
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