Exceptional service in the national interest



Solving Graph Laplacians for Complex Networks SIAM LA15, Atlanta, Oct. 2015

Erik Boman, Sandia National Labs Kevin Deweese and John R. Gilbert, UCSB





Outline



- Graph Laplacians
- Linear systems and preconditioners
- Normalization
- Empirical study with Trilinos
- Nearly-optimal combinatorial solvers
 - Kelner et al.'s simple iterative method
- Conclusions

Complex Networks: Numerical Computing



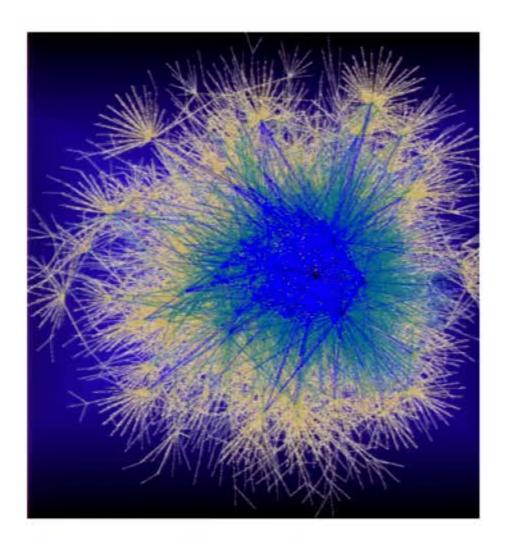
Complex networks often analyzed by

- Degree distribution
- Clustering coefficient
- Centrality metrics

Less attention on numerical linear algebra:

- Linear system: Ax=b
- Eigenproblem: Ax = λx

Well studied for PDEs, but not for complex networks.

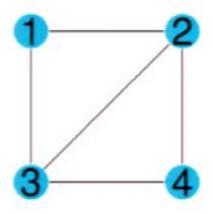


BGP graph (credit: Richardson, Chung) http://math.ucsd.edu/~fan/graphs/gallery





| Symbol | Matrix |
|-------------------------------|-------------------------------|
| Α | Adjacency matrix |
| D | Diagonal vertex degree matrix |
| L = D-A | Graph Laplacian |
| $L_N = D^{-1/2}(D-A)D^{-1/2}$ | Normalized Laplacian |



$$\begin{pmatrix} 2 & -1 & -1 & 0 \\ -1 & 3 & -1 & -1 \\ -1 & -1 & 3 & -1 \\ 0 & -1 & -1 & 2 \end{pmatrix}$$

Solving Linear Systems



Different research communities, different approaches!

- Numerical linear algebra
 - Empirical focus
 - Analysis for model problems sufficient
 - Main application: discretizations of PDEs
 - Good and robust software for solving large systems
- CS Theory
 - Focus on theory and complexity
 - Worst-case analysis
 - Main target: graph Laplacians, SDD systems
 - Software not important (some Matlab codes)
- Network Science
 - Just a tool don't care how it's done

Solvers and Preconditioners



- Sparse direct factorization only viable for small problems
- Stationary iterations (Jacobi, Gauss-Seidel) converge but quite slowly
- Conjugate gradients or Chebyshev acceleration reduces #iterations.
 - Key is to find good preconditioner M≈A
- Classic "black-box" algebraic preconditioners:
 - Jacobi (diagonal)
 - Symmetric Gauss-Seidel (SGS)
 - Incomplete Cholesky (IC)
- Algebraic multigrid (AMG)
 - Developed for PDEs on meshes, not complex networks
 - Recent progress tuning for complex networks (LAMG)

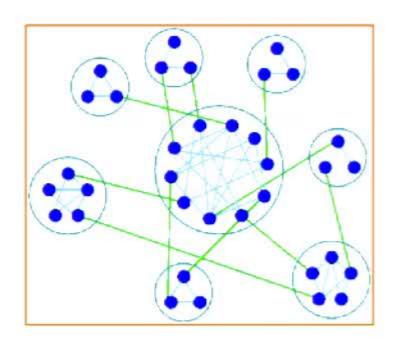
BTER Graph Generator



How to generate realistic graphs/networks?

We use BTER: Block Two-level Erdös-Renyi

- Kolda, Pinar, Seshadri (2014)
- Captures skewed degree distributions
 - Not necessarily power-law
- Has community structure
- Able to "fit" real data
 - Degree distribution
 - Clustering coefficient



Experiments



- Study two groups of graphs
 - Real networks from UF and SNAP collections
 - Social networks, web graphs, collaboration networks, etc.
 - 25 graphs, up to 735K vertices (3.5M edges)
 - Synthetic graphs (BTER)
 - Log normal degree distribution, but vary sizes and avg. degree
- Solve singular Lx=b where the solution is a random vector, using projected PCG
 - Null-space is just the constant vector
 - Use Trilinos software (next slide)
- Solvers have two phases
 - Setup (preconditioner setup or symbolic+numeric factorization)
 - Solve (CG iteration or triangular solves)

Trilinos Computational Science Toolkit



- Collection of ~60 packages
 - Heroux et al., Sandia
- Trilinos Capabilities:
 - Scalable Linear & Eigen Solvers
 - Discretizations, Meshes & Load Balancing
 - Nonlinear & Optimization Solvers
 - Software Engineering Technologies & Integration
- Parallel:
 - MPI for distributed memory
 - Growing support for sharedmemory (OpenMP, pthreads, CUDA)

Packages we used:

- Tpetra: Matrices & vectors
- Belos: Iterative solvers
- Ifpack2: Preconditioners
 - Jacobi, Gauss-Seidel
 - Incomplete factorizations
 - Subgraph preconditioners
- MueLu: Multigrid

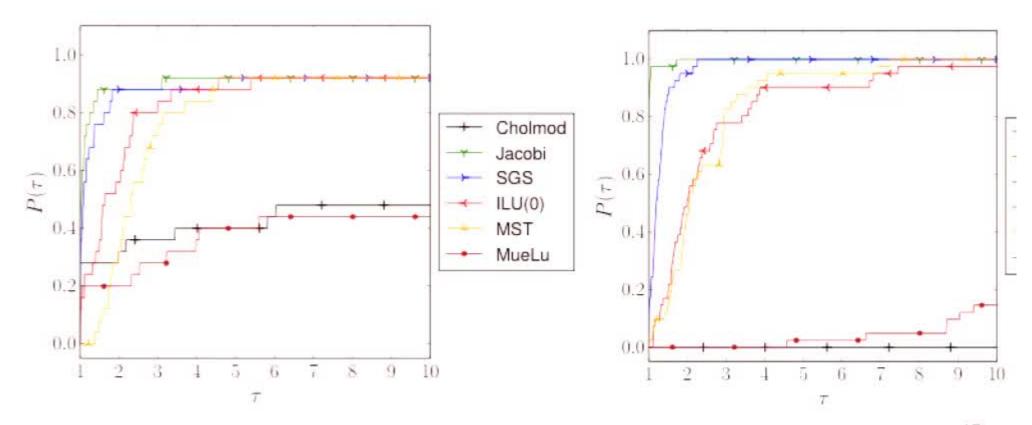
Performance Profile



Total time = Setup time + solve (iteration) time

UF real networks

BTER



Why do BTER differ from real graphs shortones

- BTER designed to match
 - Degree distribution
 - Clustering coefficient
 - Not eigenvalues!
- We tested BTER replica of Amazon-2008 network
 - Check #iterations for Laplacian solve

| Graph | Cond.no. | Jacobi | SGS | ILU(0) | AMG |
|----------|----------|--------|------|--------|-----|
| Original | 1.5e5 | 3233 | 1290 | 1211 | 216 |
| BTER | 2.0e4 | 726 | 349 | 336 | 150 |

Combinatorial Preconditioners: Great in Theory



Core Idea: Construct a sparser graph that is a good spectral approximation (spectral sparsifier), use this as preconditioner.

- Typically, use a carefully chosen subgraph
 - For example, spanning tree + "a bit more"
- First proposed by Vaidya ('90, unpublished)
 - Described and analyzed in [Bern et al. '06], implemented by [Chen and Toledo, '03]
- Support theory extensions [B., Hendrickson, '03]
- A decade of improving complexity for Laplacian/SDD solvers
 - Significant work on "near optimal solvers"
 - Spielman & Teng ('04,'05), Koutis-Miller-Peng ('10,'11), others...
 - Kelner et al. ('13): dual randomized Kaczmarz
 - Lee & Sidford ('13): coordinate descent

Are They Competitive?



- Most combinatorial near-optimal solver/preconditioners are very complicated and have never been implemented
- The recent KOSZ/DRK method is simpler:
 - Solves a dual problem on the edges of the graph
 - Corresponds to flows in an electrical network
 - Randomly sample a cycle, update flow along edges, repeat
 - This is randomized Kaczmarz (on a dual problem)
 - No CG required as convergence is provably good without
- Two recent papers evaluate this method:
 - Hoske, Lukarski, Meyerhenke, Wegner (2015)
 - B., Deweese, Gilbert (2015)
 - Both conclude KOSZ/DRK is not competitive on unweighted graphs