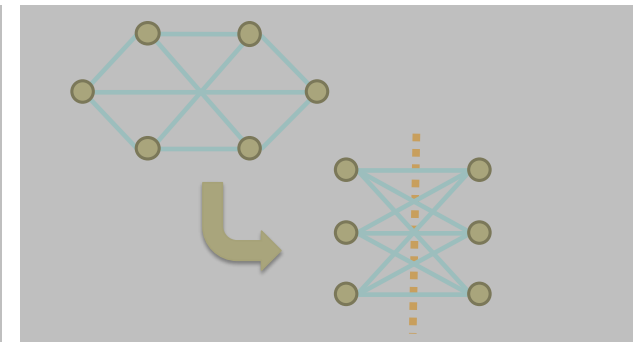
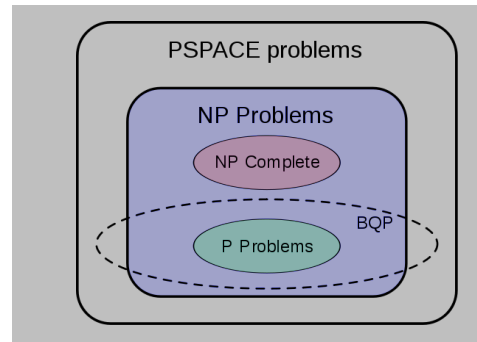
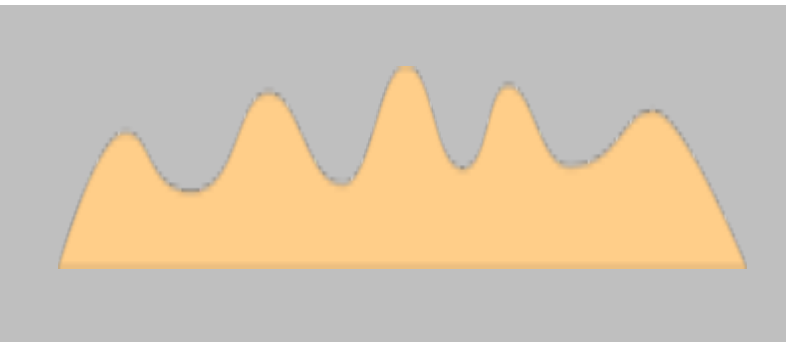


*Exceptional service in the national interest*



# Quantum Approximation Algorithms

Ojas Parekh and Ciaran Ryan-Anderson

SIAM Annual Meeting, 2017



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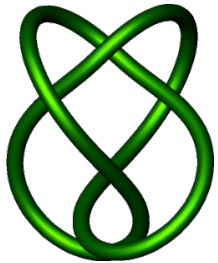
# Why quantum algorithms?

- Potential power of quantum resources is too great to ignore
- Need quantum algorithms to guide quantum hardware investment and development
- Quantum perspective has inspired new classical algorithms!
- Desire for novel quantum applications and techniques

# Limited bag of tricks for speedups

50+ algorithms: <http://math.nist.gov/quantum/zoo>

## Phase Estimation (ca. 1994)



- Factoring
- Quantum chemistry
- Linear systems
- Topological invariants

## Amplitude Amplification (ca. 1996)



- Unordered search
- Graph/network properties
- Data collision problems
- Matrix product verification

## Hamiltonian Simulation (ca. 1996)



- Quantum chemistry
- Linear systems
- Maze solving

## Quantum Walk (ca. 2002)



- Boolean formula evaluation
- Spatial search
- Quantum chemistry

New quantum algorithmic approaches are desperately needed!

# State of quantum “speedups”



- **Unproven exponential speedup:**  
Shor’s quantum factorization algorithm
- **Provable polynomial speedup:**  
Grover’s quantum search algorithm
- **Provable exponential resource advantage  
(in specialized models of computation):**  
Query and communication complexity

# Quantum bits

## State space

**Classical bit:**  
(bit)



OR



1 = Head

0 = Tail

{0, 1}

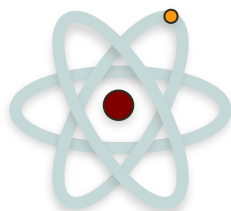
**Prob. bit:**  
(p-bit)



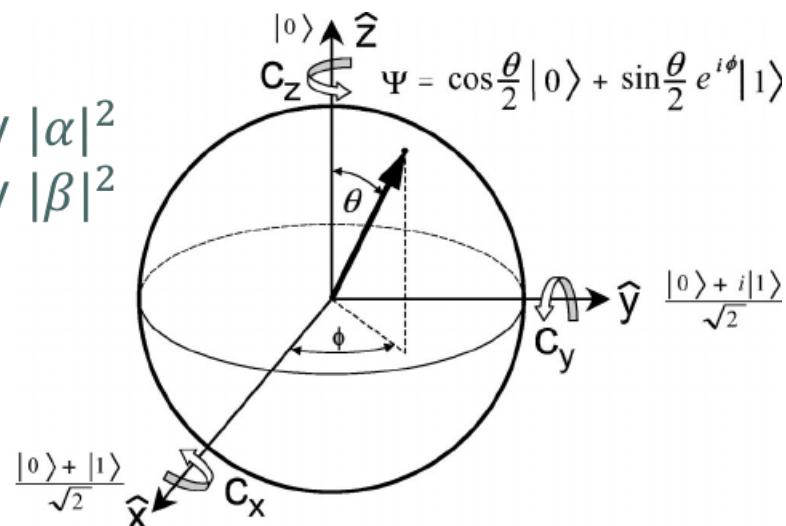
0 with probability  $1 - p$   
1 with probability  $p$



**Quantum bit:**  
(qubit)



$\alpha|0\rangle + \beta|1\rangle$   
0 with probability  $|\alpha|^2$   
1 with probability  $|\beta|^2$



# Quantum gate

Can take the “square root” of ordinary logic gates

Conventional logic gate:

NOT

yes  $\rightarrow$  no

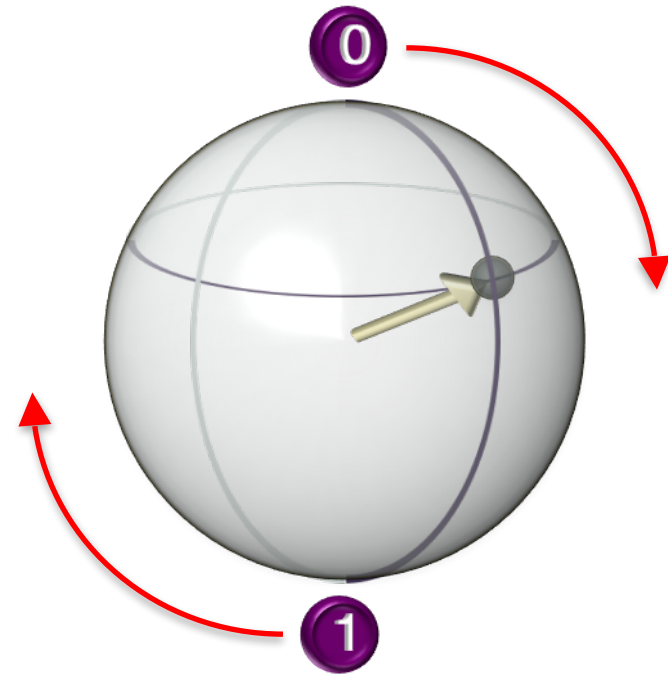
no  $\rightarrow$  yes

Quantum logic gate:

$\sqrt{\text{NOT}}$

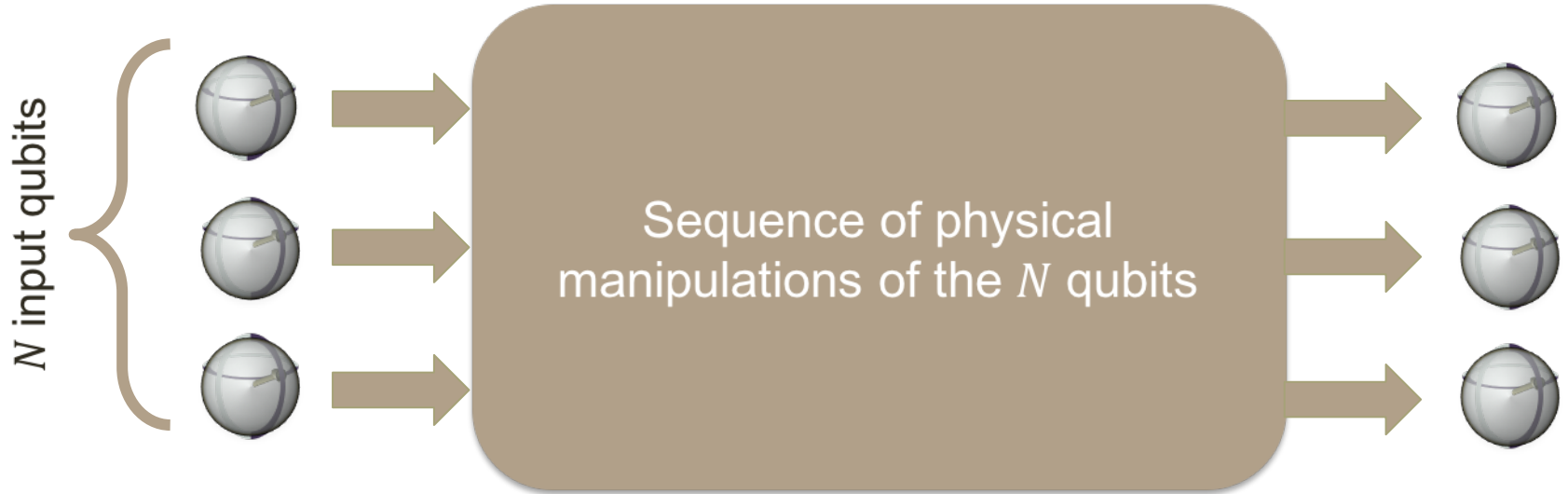
yes  $\rightarrow$  50/50 chance of yes or no

no  $\rightarrow$  50/50 chance of yes or no

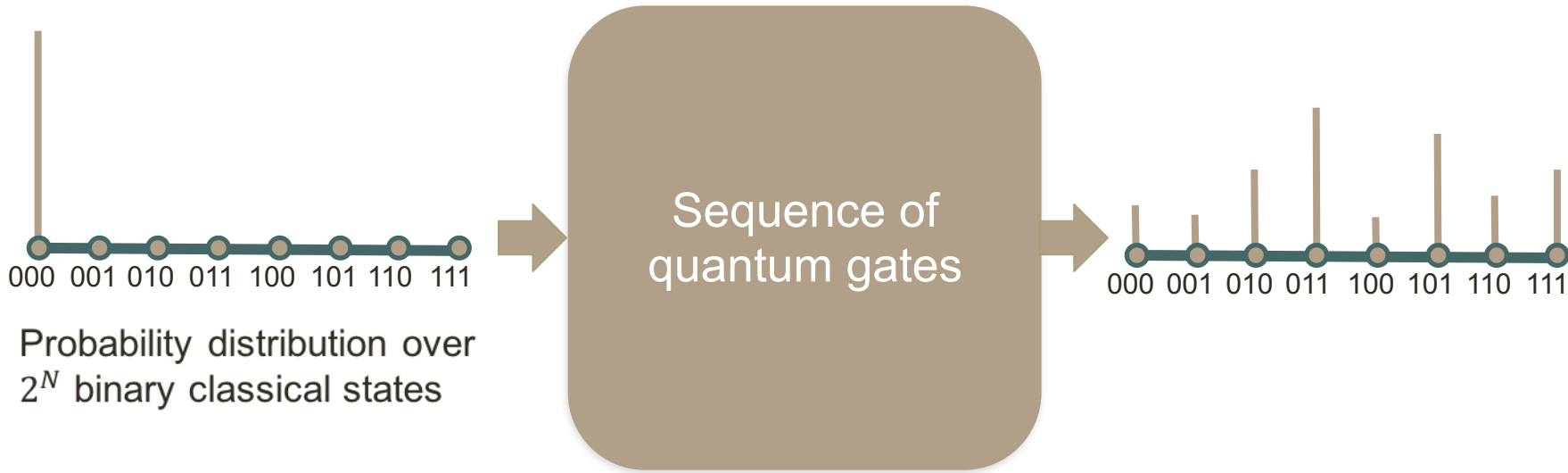


# Quantum algorithm

Physically



Conceptually



# Entanglement by analogy

Physical world



Probabilistic bits



Entangled qubits

Superposition space  
(possible measurement outcomes)



prob.  $\frac{1}{4}$



prob.  $\frac{1}{4}$



prob.  $\frac{1}{4}$



prob.  $\frac{1}{4}$



prob.  $\frac{1}{2}$



prob. 0



prob. 0



prob.  $\frac{1}{2}$

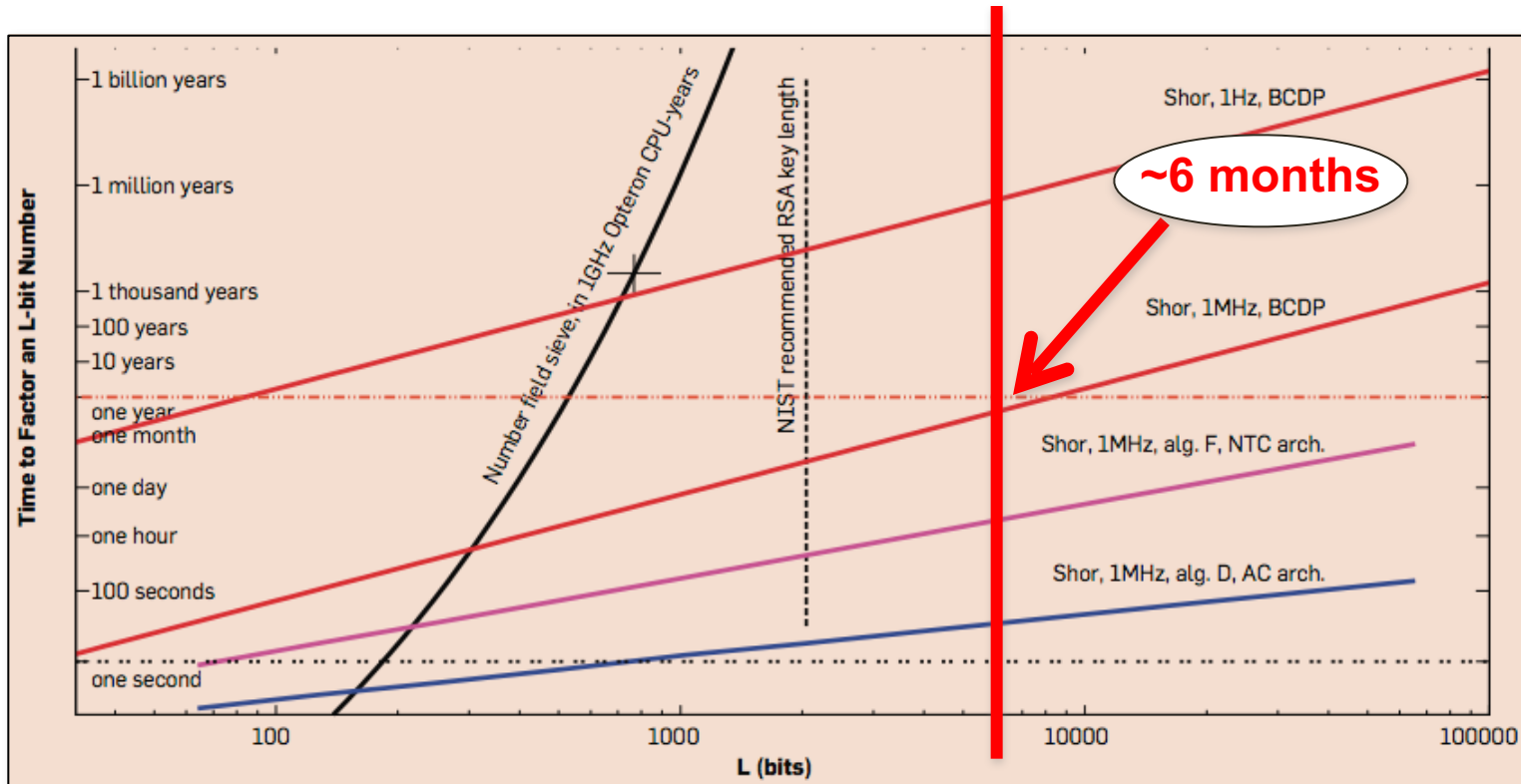
The entangled qubits will always match, even if measured at different times and across space!



# Quantum factoring

An exponential speedup

If all the silicon in the world's crust were converted to Pentium chips, it would take the age of the universe to factor a 5,000-bit number.

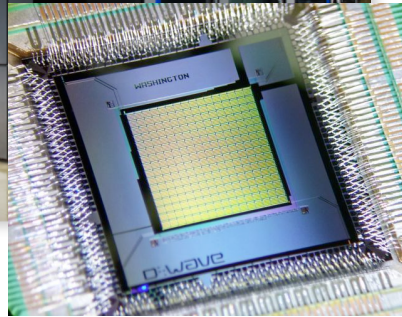


A blueprint for building a quantum computer, R. van Meter & C. Horsman, *Comm. ACM*, (2013) doi:10.1145/2494568

# Adiabatic quantum computing

Not a 'universal' computer; may have no speedup

a.k.a. "quantum annealing"



$$\min_{x \in \{0,1\}} \left( \sum_{i,j=1}^n J_{ij} x_i x_j + \sum_{i=1}^n h_i x_i \right)$$

This problem is "NP-hard:" it is unlikely that even a quantum computer could solve it efficiently.

An adiabatic quantum computer **could** be made universal, if the technology were modified to allow the qubits to interact in more interesting ways.

# Metrics status

Where we are, and where we might go

Metric	2016	2026
Universal q. computer	~10 qubits, 100 ops	~1,000 qubits, 10,000 ops
1-qubit gates	~1 in 10,000 error rate	Scalable logical qubit
2-qubit gates	~1 in 100 error rate	Scalable logical qubit
Analog q. simulator	~1,000 qubits	~10,000 qubits
Quantum annealer	~1,000 qubits	~10,000 qubits

**Benchmark:** 50 qubits is beyond the simulation capabilities of today's best supercomputers.

# Testbed QCs

**Google:** 49-qubit goal by December 2017.

**NSF:** \$3M/yr Ideas Lab: Practical Fully-Connected Quantum Computer Challenge (PFCQC), November 2017

**DOE:** \$5M/yr Quantum Testbed User Facility (pending Congressional budget action)

**IBM:** Open-Access "Quantum Experience" online since 5/16: 40k users, 270k experiments, 15 published papers

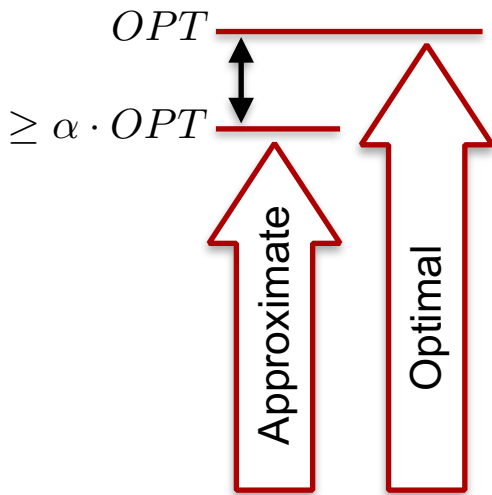
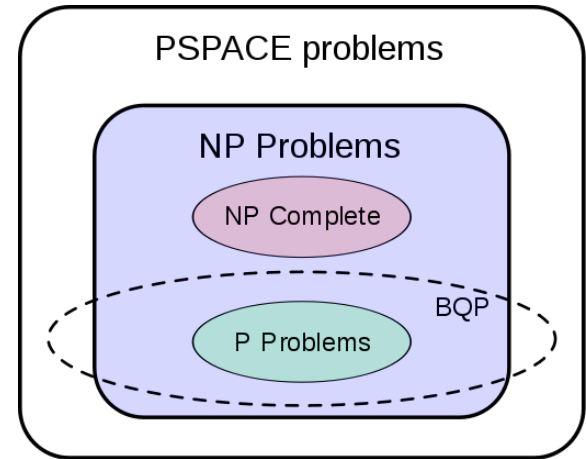
**ETH Zurich:** 45 *simulated* qubits on Cori II (#5) (4/17) [1]



[1] Haener and Steiger, arXiv:1704.01127 (2017).

# Quantum Approximation Algorithms

**Motivation:** hard to efficiently find optimal solutions for NP-complete optimization problems, **even for quantum computers**



**Approach:** an *approximation algorithm* efficiently produces a near-optimal solution with a mathematically provable bound on quality

**Benefit:** *quantum approximation algorithms (QAA)* direct quantum resources towards **higher-quality solutions** instead of faster **running times**, sidestepping barriers to quantum speedups

# The QAOA

The **Q**uantum **A**pproximate **O**ptimization **A**lgorithm was introduced by Farhi et al. in 2014

$$e^{i \sum_i \beta X_i} e^{i\gamma \sum_{ij \in E} Z_i Z_j} |+\rangle^{\otimes n}$$

## **Only known quantum approximation algorithm**

Classical approximation algorithms have been studied since the 1960s

- Can be viewed as a discretization of adiabatic quantum computing
- Results in low-depth quantum circuits
- Generic framework for combinatorial optimization problems

# QAOA for Max 3-XORSAT

Goal of Max 3-XORSAT is to satisfy max number out of  $m$  given clauses:

$$(x_1 \oplus x_3 \oplus \neg x_4), (\neg x_1 \oplus x_2 \oplus x_3), \dots$$

Restricted version: each variable appears in at most  $d$  clauses

**Farhi et al. showed that QAOA beat the best known classical approx alg:**

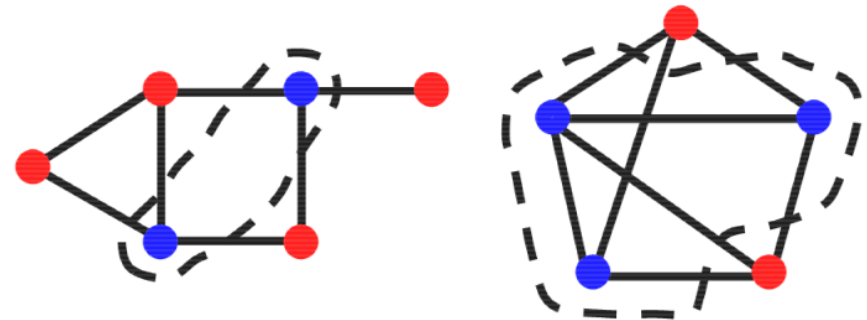
<i>Authors</i>	<i>Year</i>	<i>Result</i>	<i>Type</i>
<i>Trevisan</i>	<i>2000</i>	$\left(\frac{1}{2} + \frac{O(1)}{d}\right) m$	<i>Classical</i>
<i>Farhi et al.</i>	<i>2014</i>	$\left(\frac{1}{2} + \frac{O(1)}{d^{3/4}}\right) m$	<i>Quantum</i>
<i>Barak et al.</i>	<i>2015</i>	$\left(\frac{1}{2} + \frac{O(1)}{\sqrt{d}}\right) m$	<i>Classical</i>
<i>Farhi et al.</i>	<i>2015</i>	$\left(\frac{1}{2} + \frac{O(1)}{\log d \sqrt{d}}\right) m$	<i>Quantum</i>

Barak et al.'s result is best possible up to constants unless P=NP

# QAOA for Maximum Cut

We show that QAOA outperforms best classical algorithm for the well-known Maximum Cut problem on  $d$ -regular triangle-free graphs with  $m$  edges

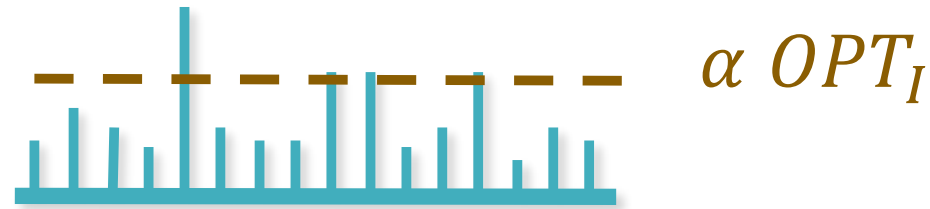
<i>Authors</i>	<i>Year</i>	<i>Result</i>	<i>Type</i>
<i>Shearer</i>	<b>1992</b>	$\left(\frac{1}{2} + \frac{0.177}{\sqrt{d}}\right)^m$	<b>Classical</b>
<i>Hirvonen et al.</i>	<b>2014</b>	$\left(\frac{1}{2} + \frac{0.281}{\sqrt{d}}\right)^m$	<b>Classical</b>
<i>Parekh et al.</i>	<b>2017</b>	$\left(\frac{1}{2} + \frac{0.303}{\sqrt{d}}\right)^m$	<b>Quantum</b>



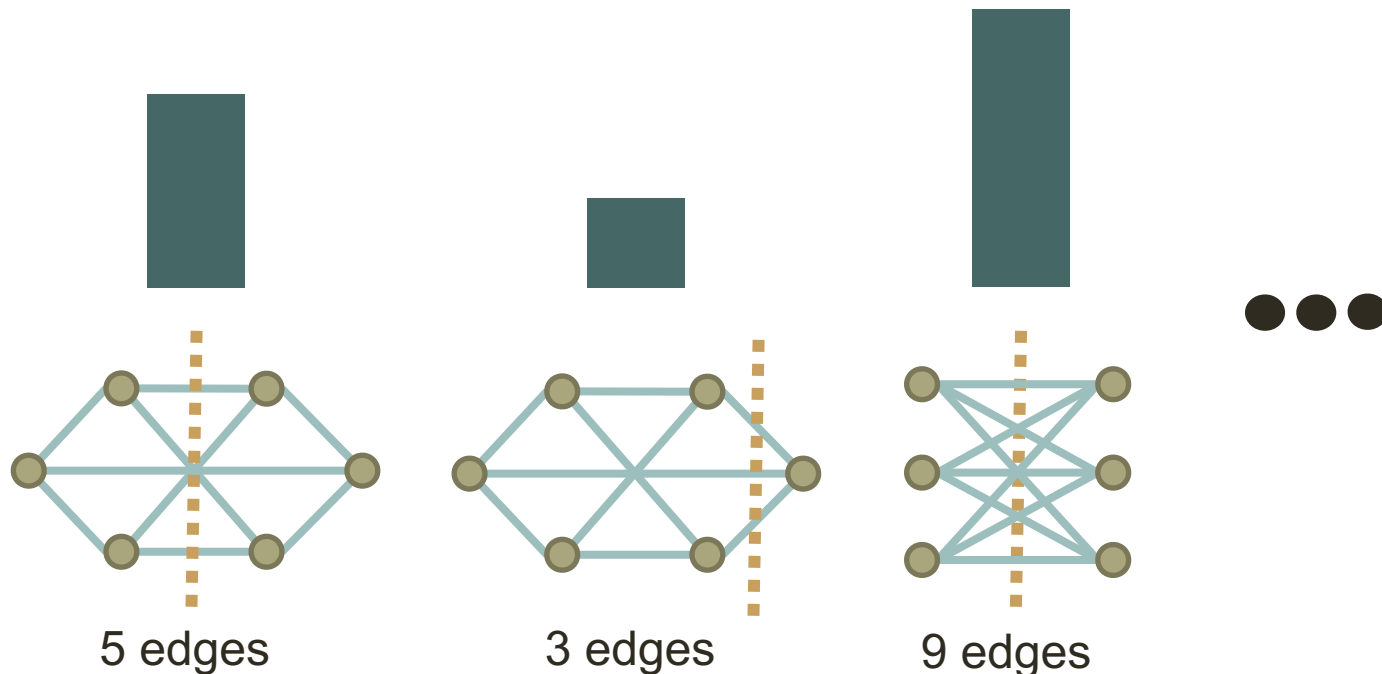
**Only known quantum approximation algorithm outperforming the best-known classical algorithm**



# Sampling vs Optimization



Our quantum algorithm allows sampling from a probability distribution on cuts in a graph, likely to yield a cut with many edges

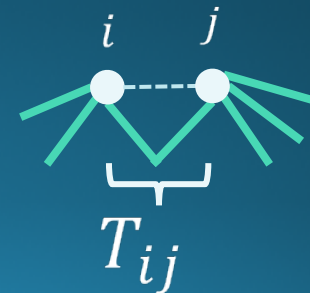
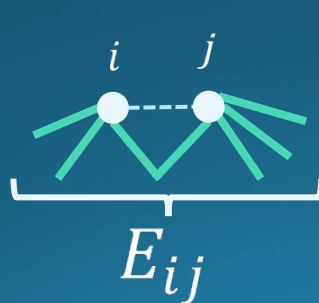
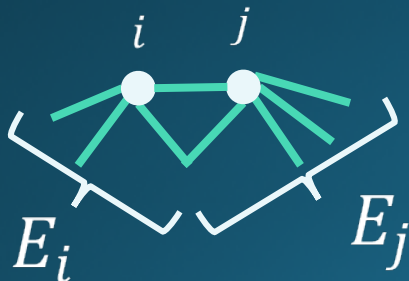


# QAOA Analysis

$$\langle C \rangle = \sum_{\langle i,j \rangle \in E} \langle C_{\langle i,j \rangle} \rangle$$

$$\langle C_{\langle i,j \rangle} \rangle =$$

$$\frac{1}{2} \left[ 1 - \frac{1}{2} \sin(4\beta) \sin(2\gamma) \left\{ \cos^{|E_i|-1}(2\gamma) + \cos^{|E_j|-1}(2\gamma) \right\} - \sin^2(2\beta) \cos^{|E_{ij}|-2|T_{ij}|}(2\gamma) \left\{ \frac{1 - \cos^{|T_{ij}|}(4\gamma)}{2} \right\} \right]$$

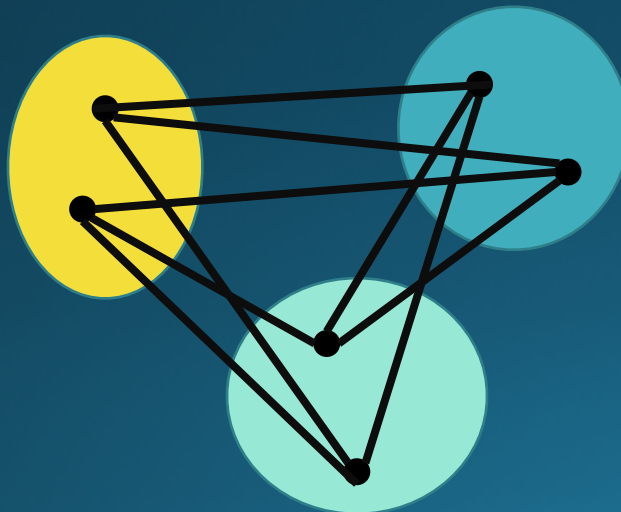


# Classical Outperforms QAOA

- Linear time algorithm:

W. Staton, *Ars Combinatoria* **10** (1980), 103-106.

Any 3-regular, connected graph (other than  $K_4$ ):



Staton:

$$C(z) \geq \frac{7}{9}m = 0.\bar{7}m$$

QAOA-1:

$$\langle C \rangle \leq 0.692451m$$